A Parallel Fast Algorithm for Applying Fourier Integral Operators

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Outline

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Fourier Integral Operators (FIOs)

The general form is

$$(Af)(x) = \int_{\mathbb{R}^d} a(x,k) e^{i\Phi(x,k)} \hat{f}(k) dk,$$

where $\hat{f}$ is the Fourier transform of f.

Fourier Integral Operators (FIOs)

Our discrete form is

$$(Af)(x_i) = \sum_{k_j \in \mathbb{R}^d} a(x_i, k_j) e^{i\Phi(x_i, k_j)} \hat{f}(k_j),$$

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Basic properties[1]:

- $a(x, k)$ is complex-valued. It is the *symbol* or *amplitude function*.
- $\Phi(x, k)$ is real-valued and is called the *phase function*.

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Convenient definitions:

- When $a(x, k) \equiv 1$, an FIO reduces to an *Egorov operator*.
- Define A_R such that $A = A_R \circ \mathcal{F}$, where $\mathcal{F}$ is the Fourier transform. (We also have $A^* = \mathcal{F}^* \circ A_R^* = \mathcal{F}^{-1} \circ A_R^*$.)
- We will simply refer to A_R as a *reduced FIO* (RFIO) since it will be our main focus.

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The discrete RFIO is then simply

$$(A_R g)(x_i) = \sum_{k_j \in \mathbb{R}^d} a(x_i, k_j) e^{i\Phi(x_i,k_j)} g(k_j).$$

Notice that the continuous Fourier transform is an RFIO where $a(x,k) \equiv (2\pi)^{-d/2}$ and $\Phi(x,k) = -2\pi x \cdot k$.

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ϵ-separation Rank

Suppose $a : \Omega_x \times \Omega_k \to \mathbb{C}$ and $\Phi : \Omega_x \times \Omega_k \to \mathbb{R}$ are smooth. Then, for a given $\epsilon > 0$, there exists $d_\epsilon > 0$ and $r_\epsilon \in \mathbb{N}$ such that, for any $X \subset \Omega_x$ and $K \subset \Omega_k$ satisfying $w(X)w(K) \leq d_\epsilon$,

$$\left| a(x,k)e^{i\Phi(x,k)} - \sum_{t=0}^{r_\epsilon - 1} \alpha_t^{XK}(x)\beta_t^{XK}(k) \right| \leq \epsilon, \quad \forall\, x \in X,\ \forall\, k \in K.$$

Summation over each $g(k)$ yields

$$\left| (A_R g_K)(x) - \sum_{t=0}^{r_\epsilon - 1} \alpha_t^{XK}(x) \left(\sum_{k \in K} \beta_t^{XK}(k) g(k) \right) \right| \leq \left(\sum_{k \in K} |g(k)| \right) \epsilon,$$

which can be rewritten as

$$(A_R g_K)(x) \approx \sum_{t=0}^{r_\epsilon - 1} \alpha_t^{XK}(x) \delta_t^{XK}, \quad \forall\, x \in X,\ \forall\, k \in K.$$

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- $w(X)w(K) \leq d_\epsilon$ only constrains product of diameters of X and K.
- $X = \Omega_x$ and $K = \Omega_k$ are both allowed, just not at once.
- If we could cover Ω_x with many low-rank approx's using $K = \Omega_k$, we could still cheaply evaluate $(A_R g)(x)$.
- But if $K = \Omega_k$, forming $\sum_{k\in K} \beta_t^{XK}(k) g(k)$ is expensive and must be done for each $X \subsetneq \Omega_X$.
- If $X = \Omega_x$, so that K is small, forming each $\sum_{k\in K} \beta_t^{XK}(k) g(k)$ is cheap!
- Then transform many (Ω_x, K) interactions into many (X, Ω_k) interactions...

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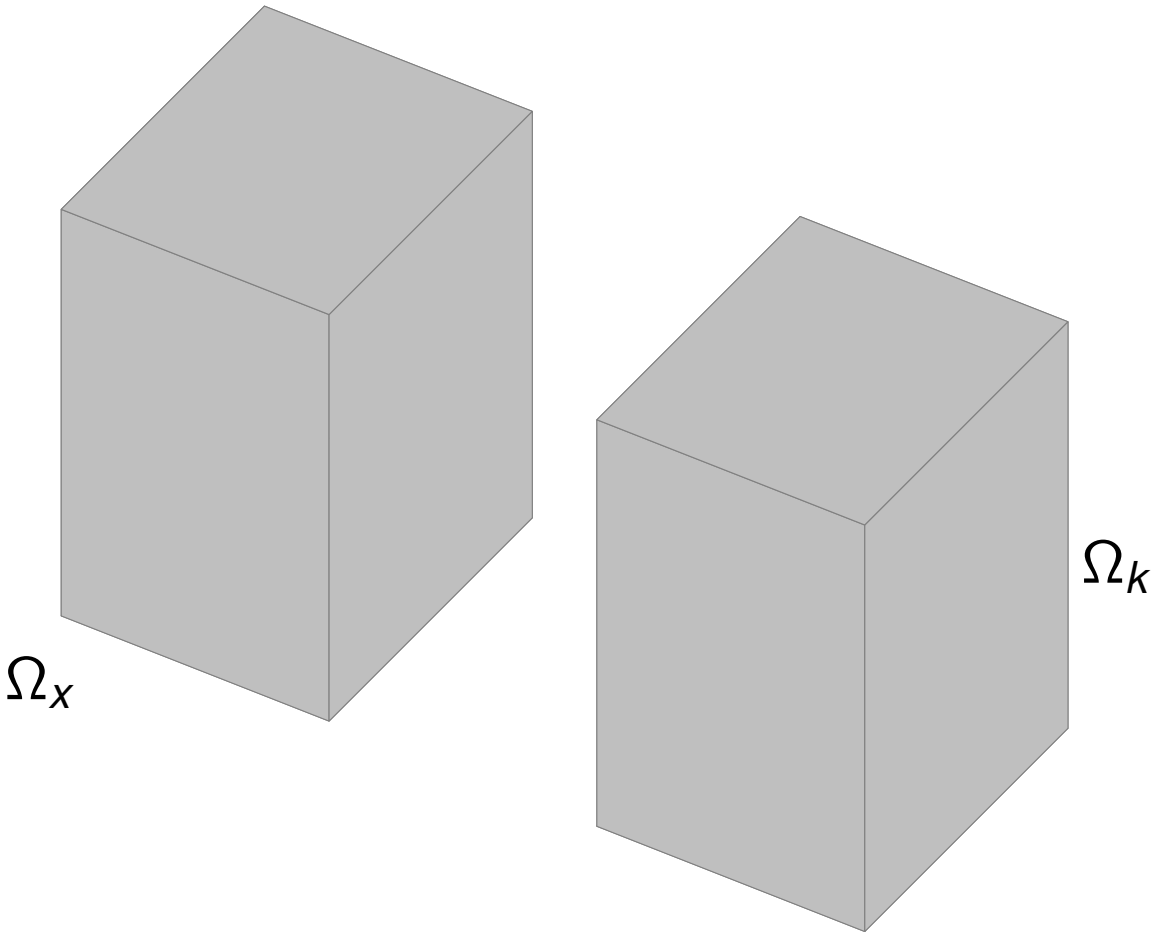
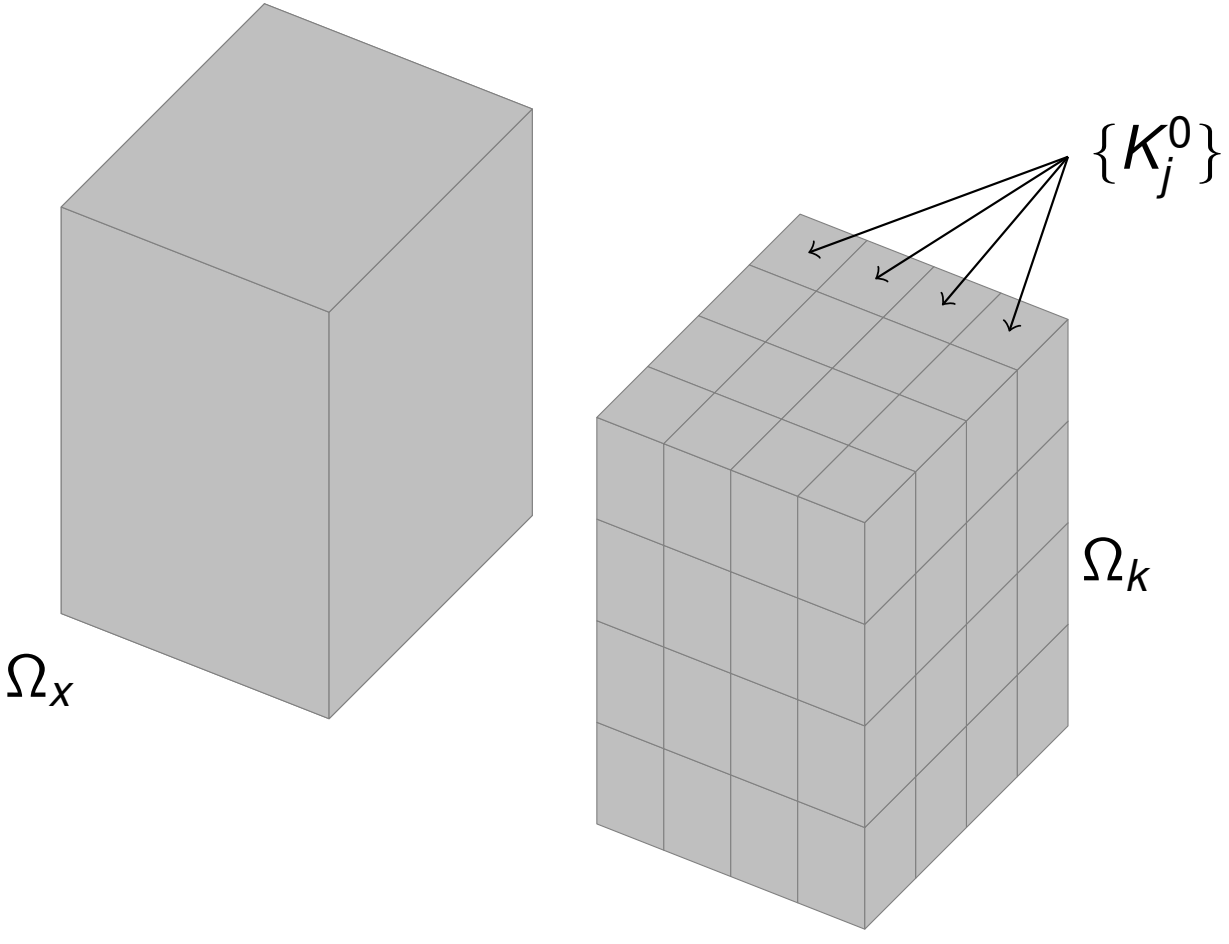
Butterfly Algorithm Example: Level 0

Choose $N = 2^m$, $m \in \mathbb{N}$, such that $w(\Omega_x)w(\Omega_k)/N \leq d_\epsilon$.



Butterfly Algorithm Example: Level 0

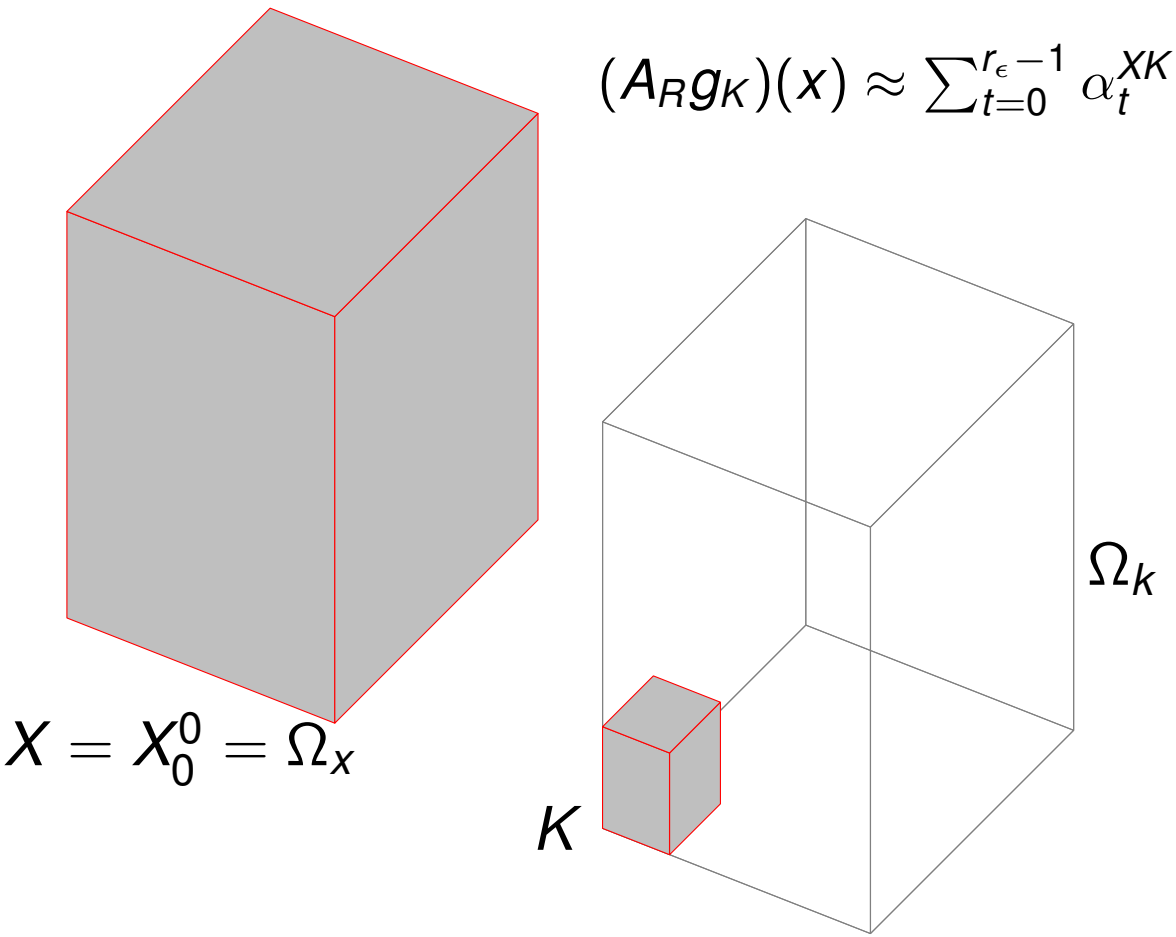
Partition Ω_k into N pieces in each dimension (i.e., $N = 4$), $\{K_j^0\}_{j=0}^{N^d-1}$.



Butterfly Algorithm Example: Level 0

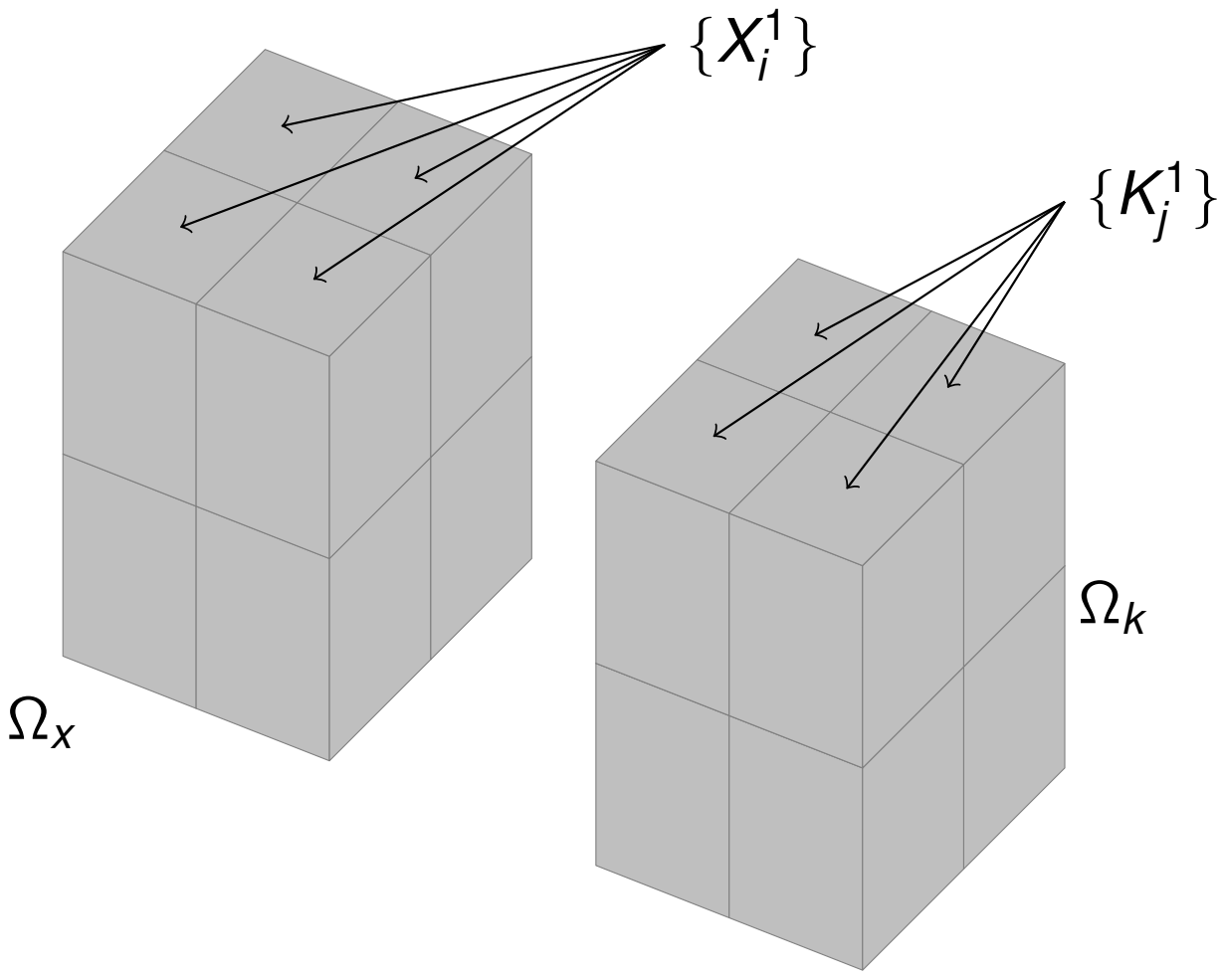
For each $K = K_j^0$, initialize $\{\delta_t^{XK}\}$ from sources in K.

$$(A_R g_K)(x) \approx \sum_{t=0}^{r_\epsilon - 1} \alpha_t^{XK}(x)\delta_t^{XK}, \ \forall x \in X$$



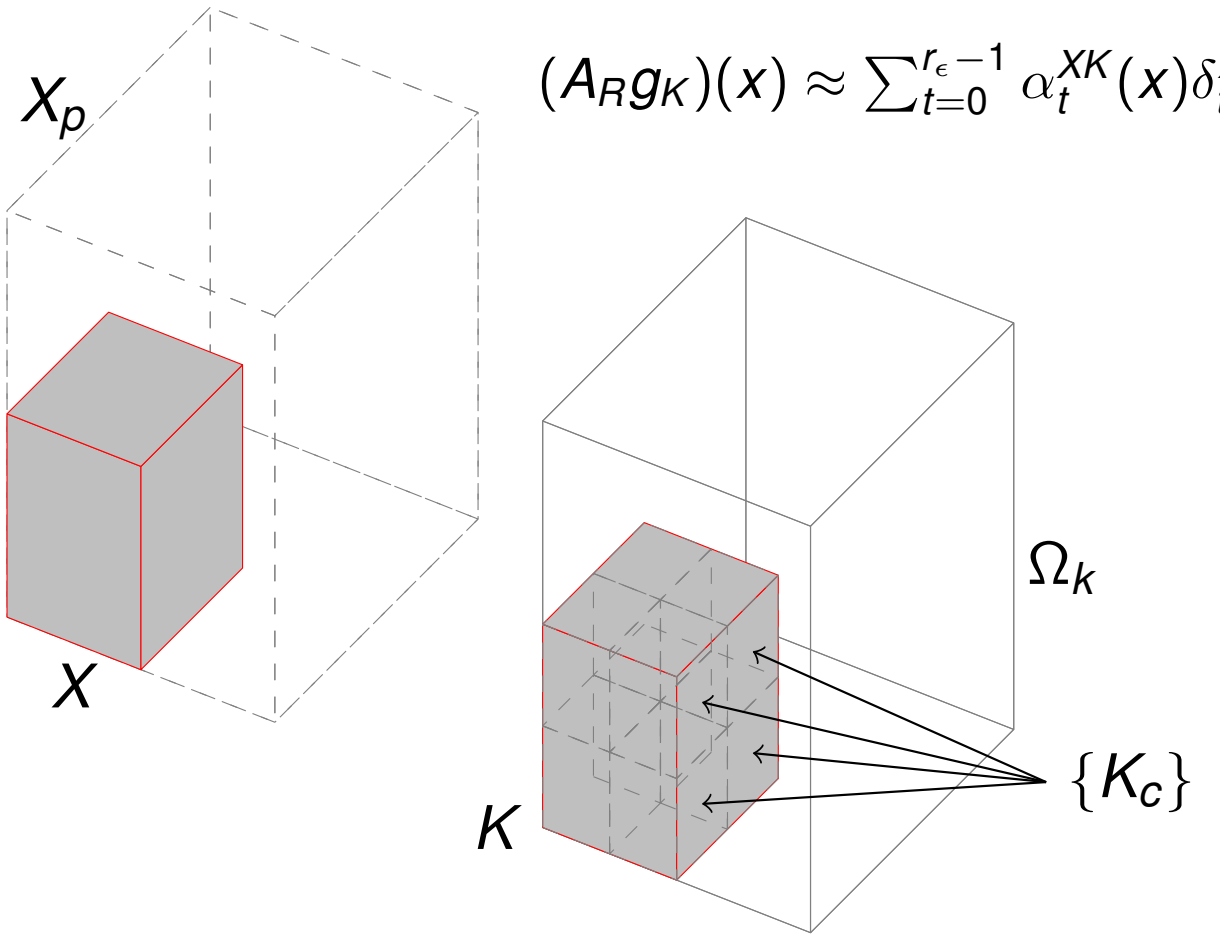
Butterfly Algorithm Example: Level 1

Refine Ω_x and coarsen Ω_k by factor of 2 to form $\{X_i^1\}$ and $\{K_j^1\}$.



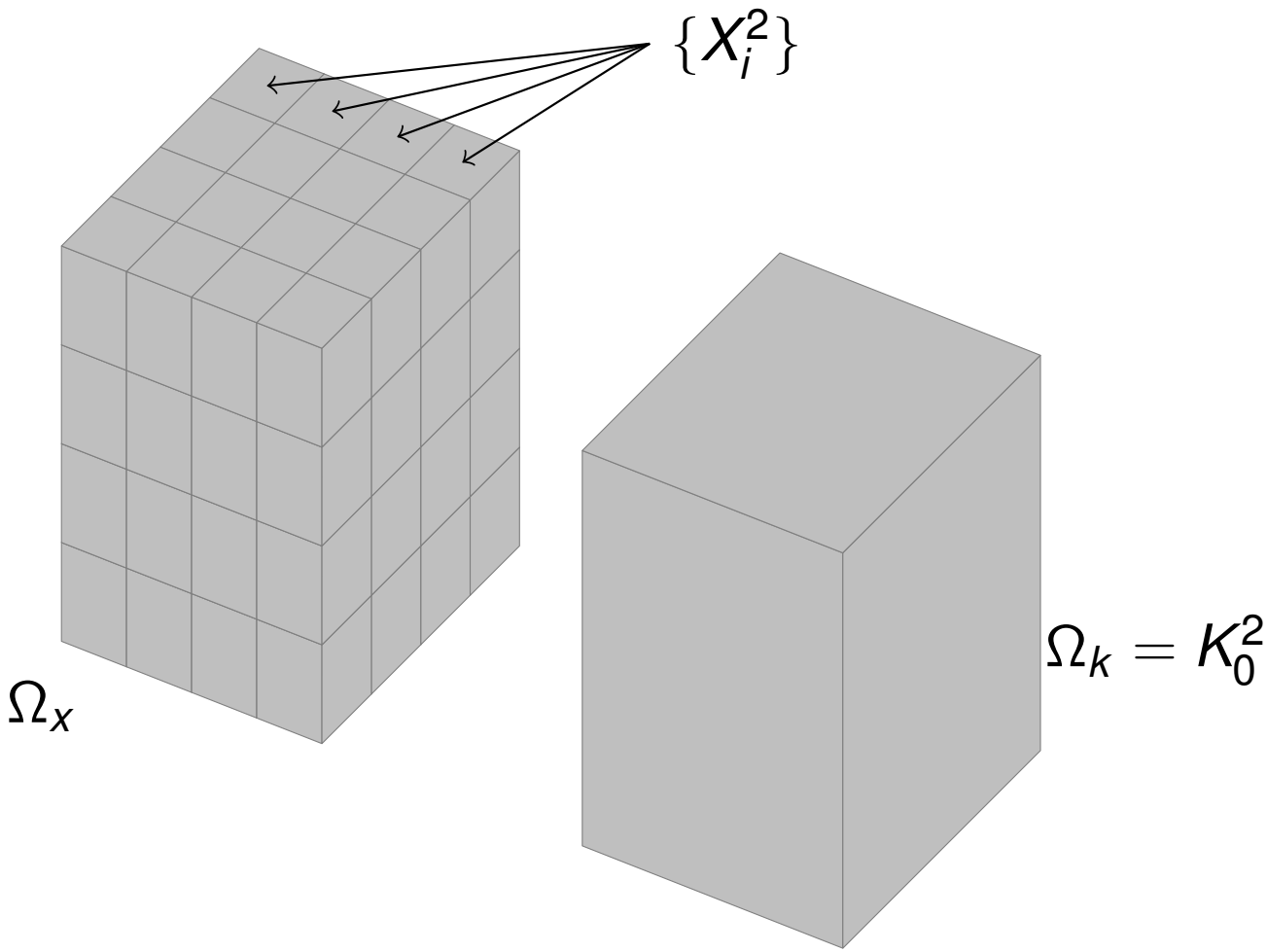
Butterfly Algorithm Example: Level 1

For each $X = X_i^1$ and $K = K_j^1$, form $\{\delta_t^{XK}\}$ by combining the weights for the interaction of K's 2^d children with X's parent.



Butterfly Algorithm Example: Level 2

Refine Ω_x and coarsen Ω_k by factor of 2 to form $\{X_i^2\}$ and $\{K_j^2\}$.



Butterfly Algorithm Example: Level 2

For each $X = X_i^2$, form $\{\delta_t^{XK}\}$ by combining the weights for the interaction of K's 2^d children with X's parent.

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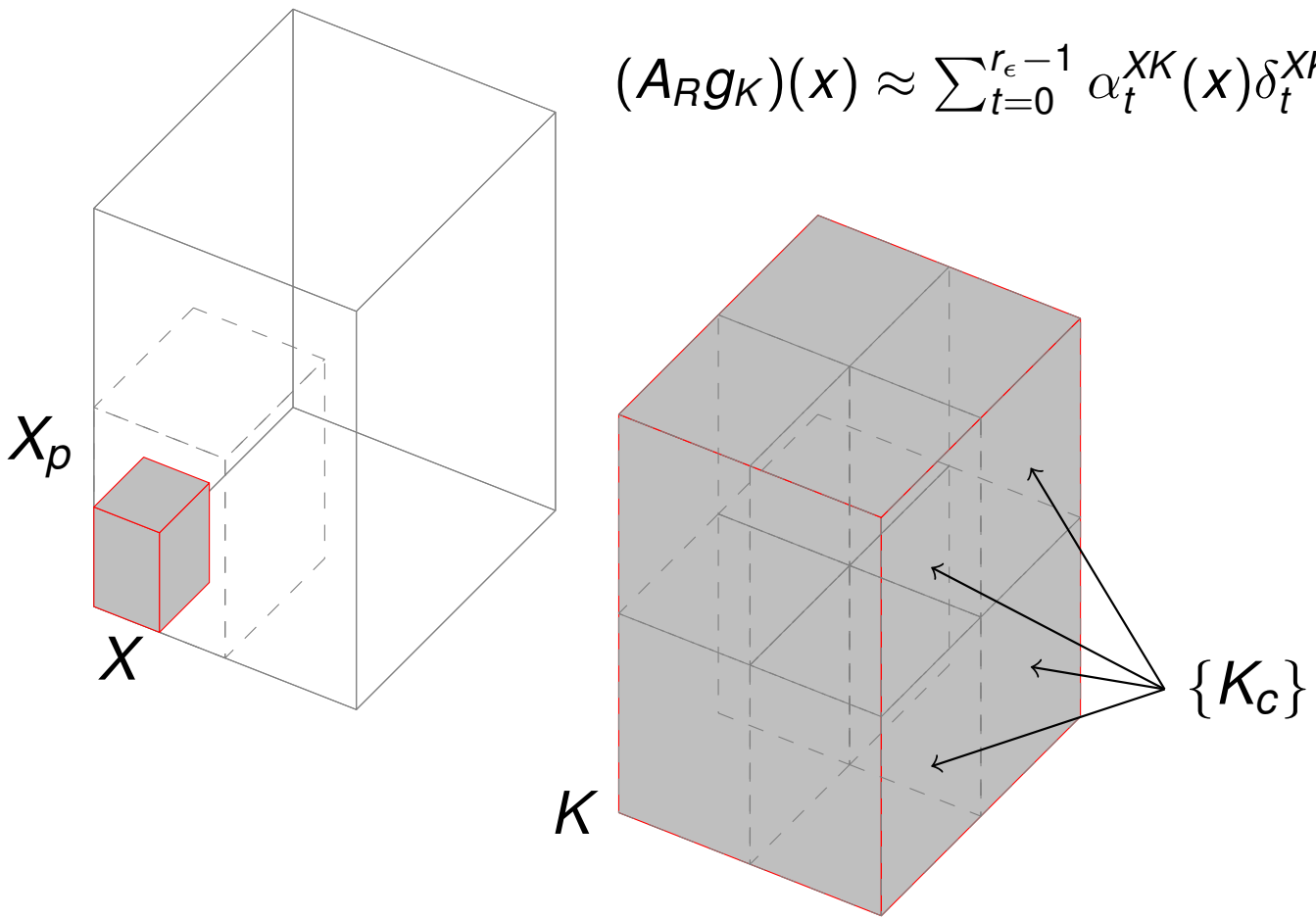

Butterfly Algorithm Example: Summary

- After completion, $K = \Omega_k$, so $g_K = g$ and we have

$$(A_R g)(x) \approx \sum_{t=0}^{r_\epsilon - 1} \alpha_t^{XK}(x) \delta_t^{XK}, \ \forall x \in X.$$

- In general, there are $\log_2(N)$ levels.
- Since $\{X_i^{\log_2(N)}\}$ covers Ω_x, one can evaluate $(A_R g)(x)$ for any $x \in \Omega_x$ with $\mathcal{O}(1)$ work.
- Each level involves N^d interactions requiring $\mathcal{O}(1)$ work each.
- Complexity is $\mathcal{O}(N^d \log_2(N))$, vs. $\mathcal{O}(N^{2d})$ naïve complexity.

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Parallelization Strategy

- Grids change at every level, but the number of interactions is constant (N^d).
- The idea is to keep the interactions (and therefore data) evenly distributed.
- Begin with Ω_k fully distributed and end with Ω_x fully distributed.
- Use power of two number of processes so that repeated bisections can be used to distribute the domains.
- Gradually migrate the bisections from Ω_k to Ω_x during the coarsening of Ω_k.
- Limit to at most N^d processes (one interaction per process).

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- Gradually migrate the bisections from Ω_k to Ω_x during the coarsening of Ω_k.
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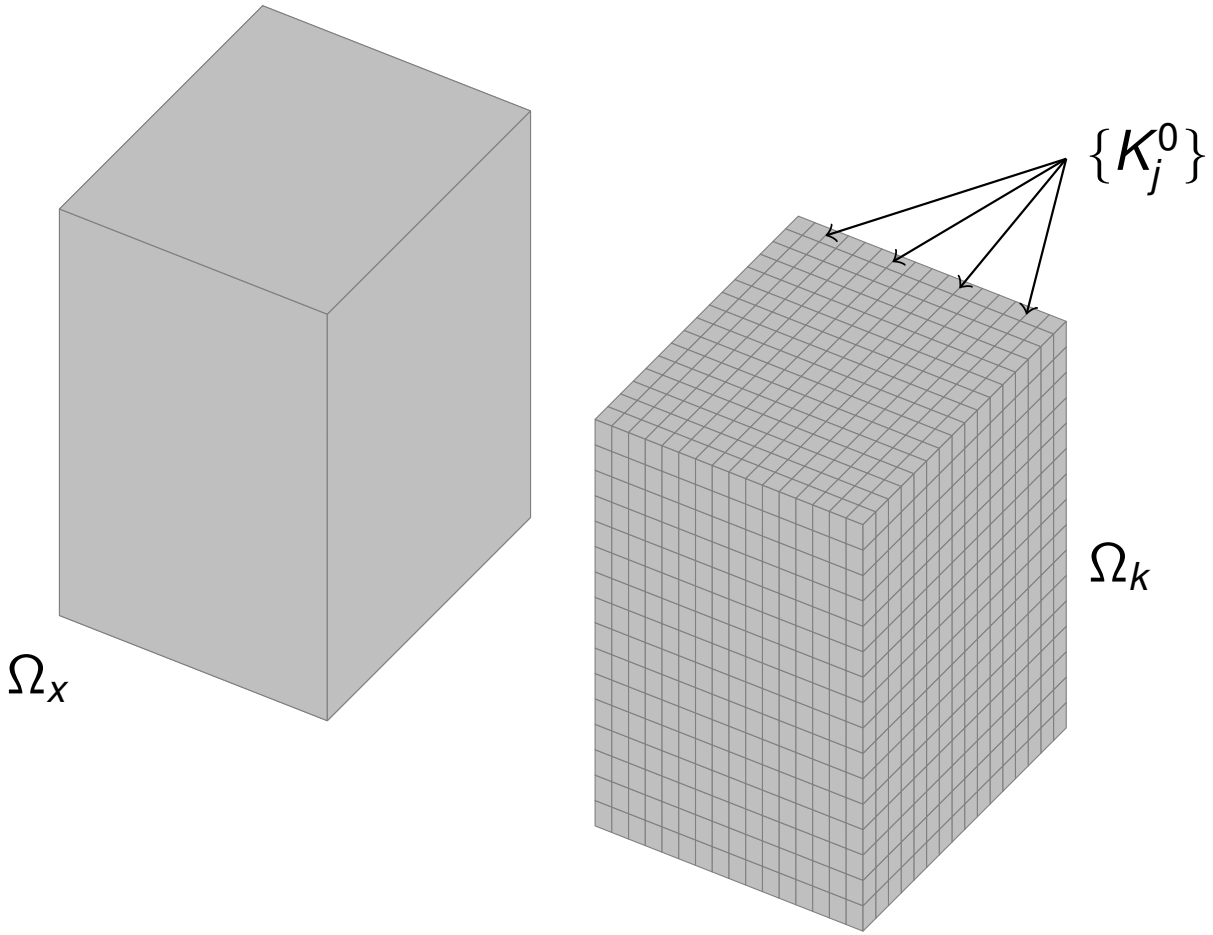
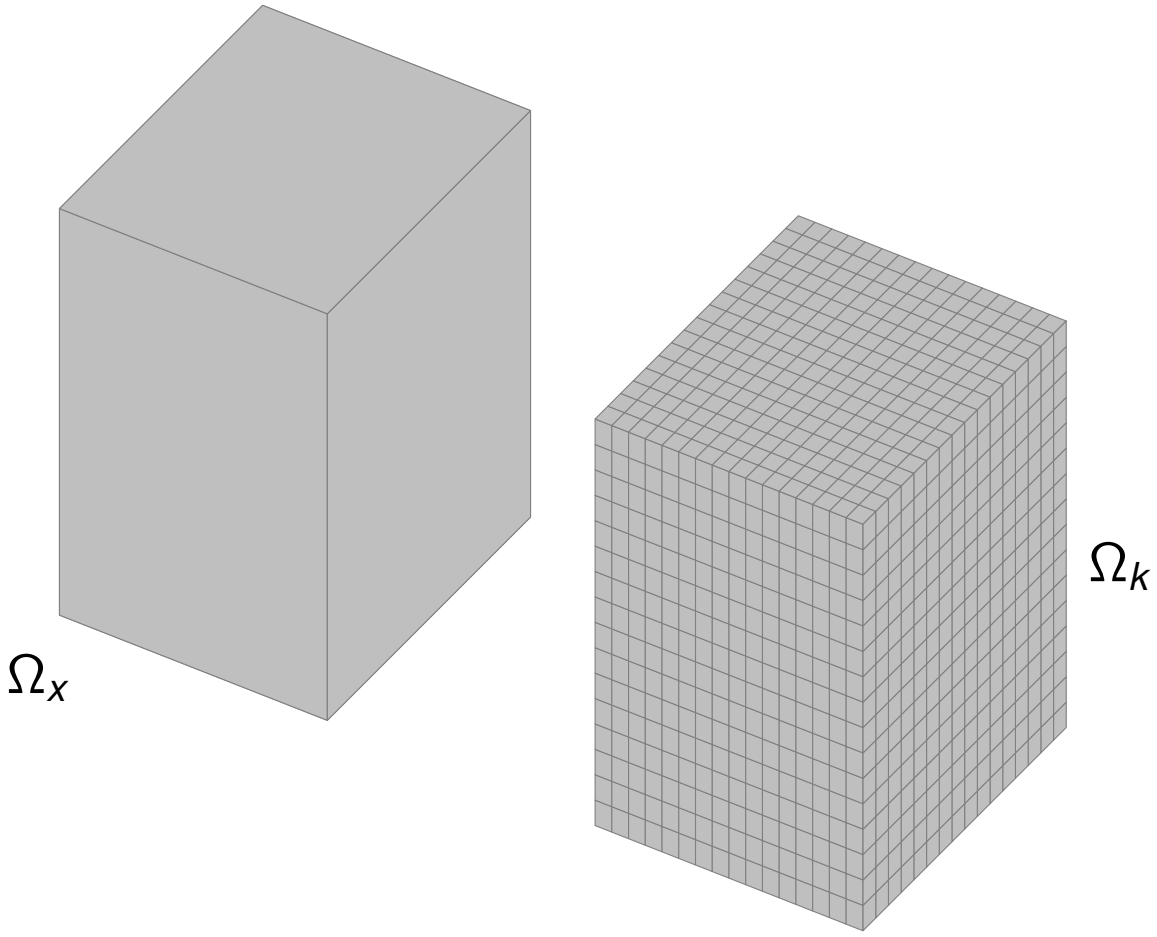
Example for $d = 3$, $N = 16$, and $p = 16$: Level 0

We begin with Ω_k partitioned into N^d pieces, $\{K_j^0\}$.



Example for $d = 3$, $N = 16$, and $p = 16$: Level 0

We cyclically bisect Ω_k $\log_2(p)$ times over the d dimensions to distribute Ω_k.



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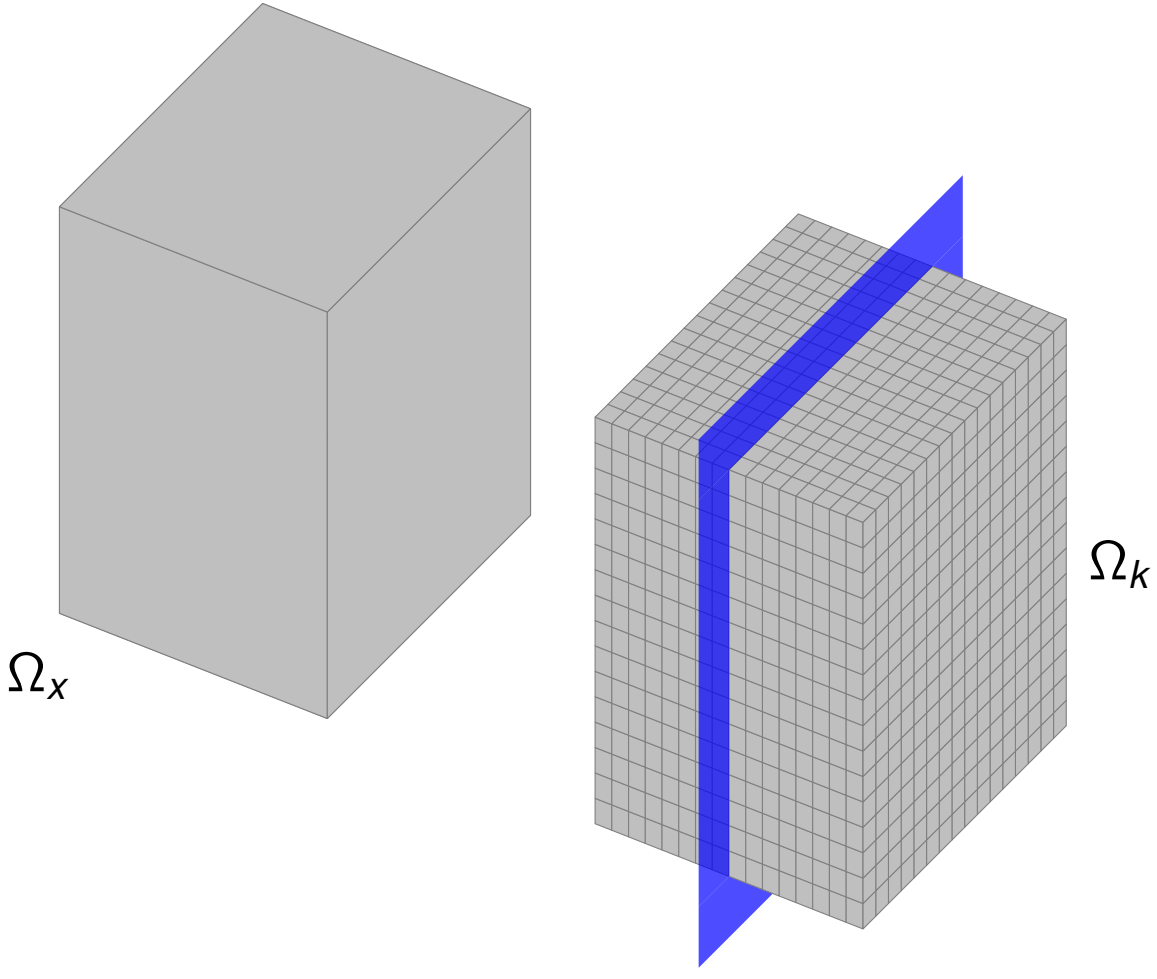

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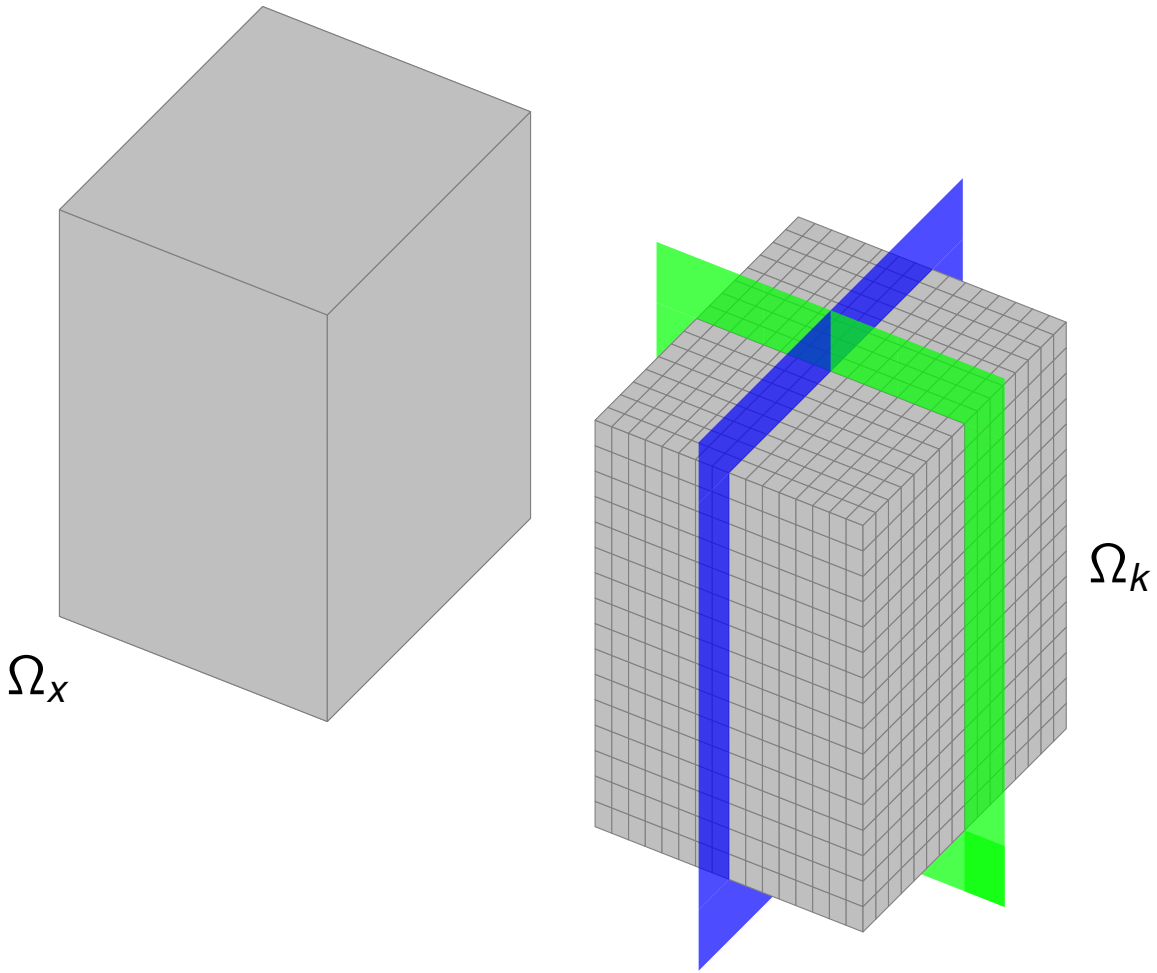

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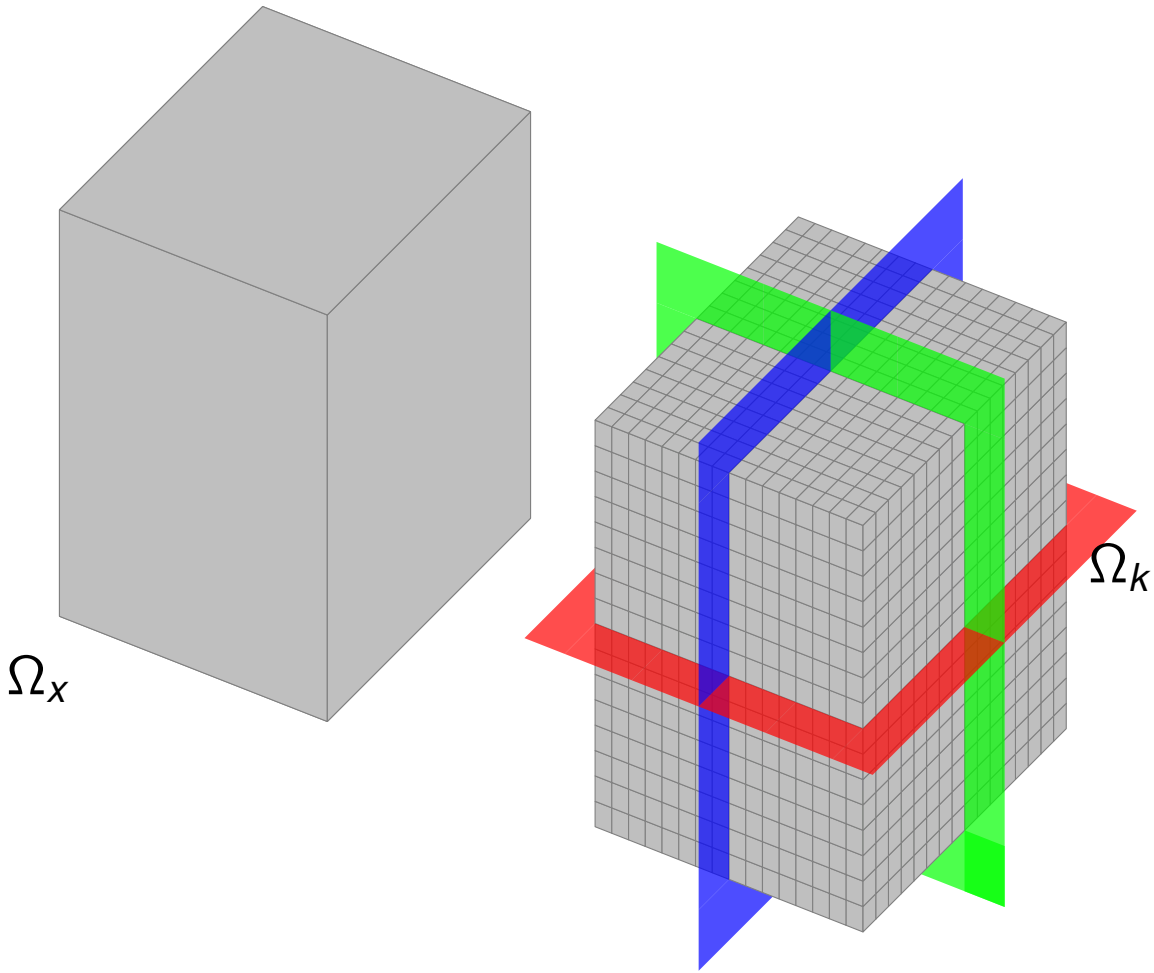

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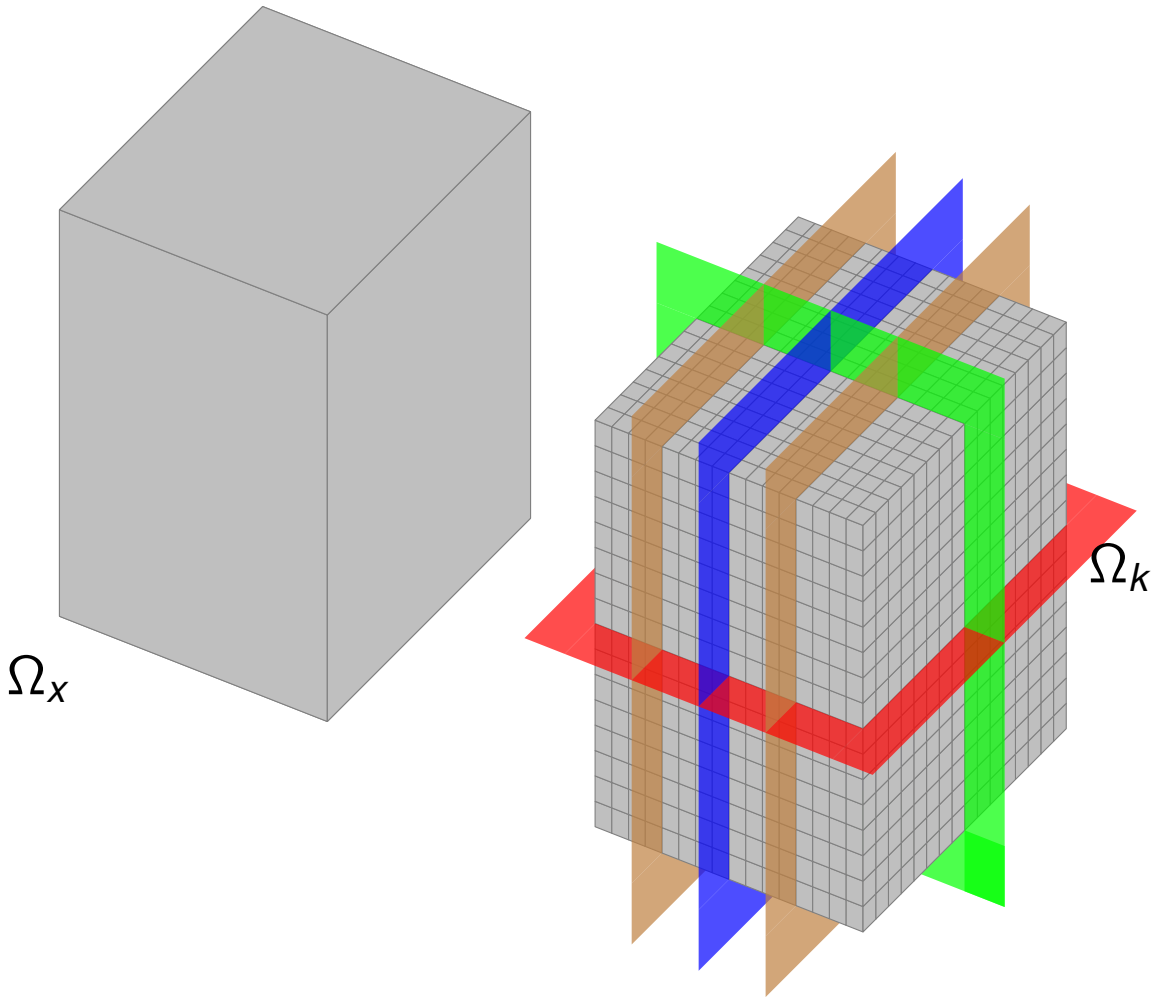
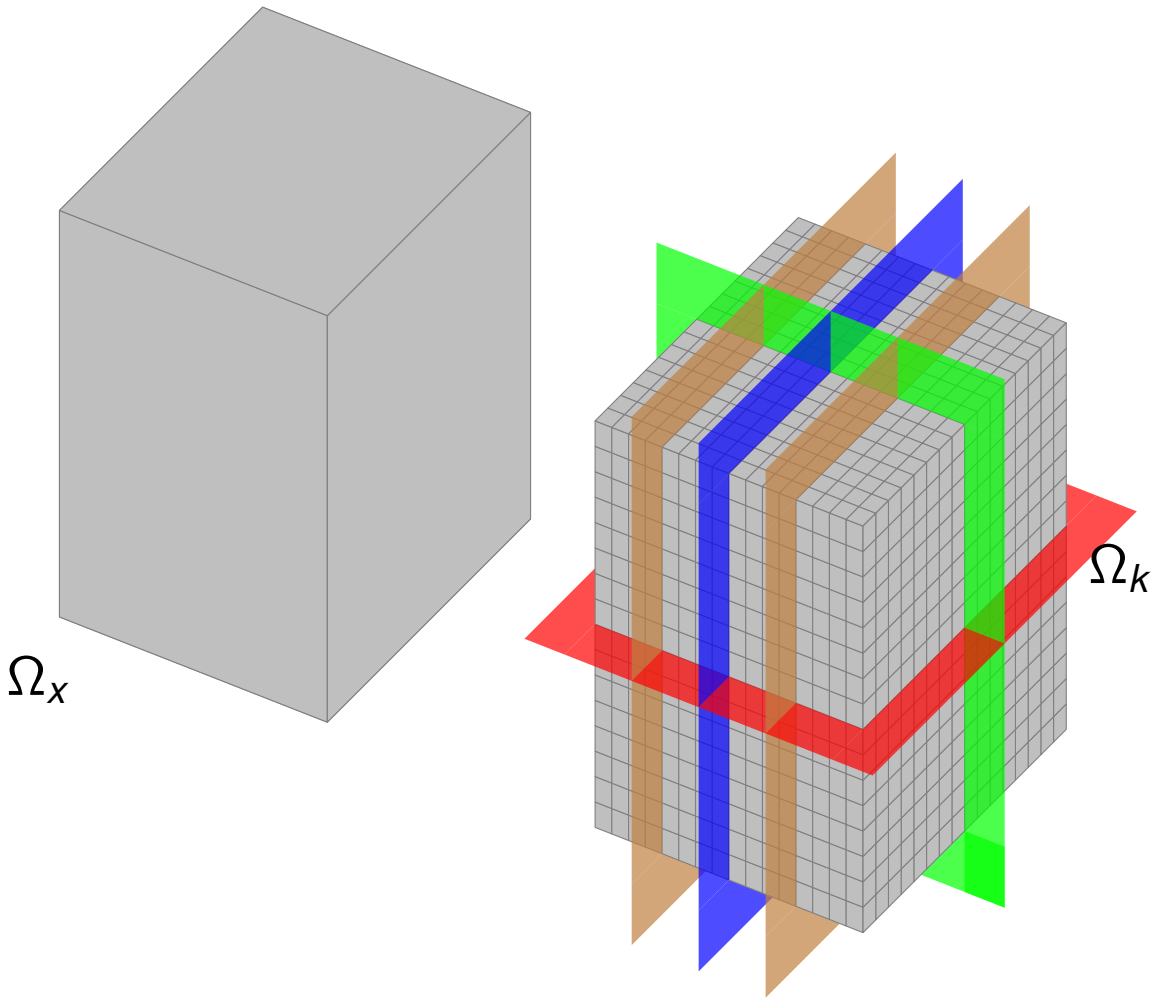


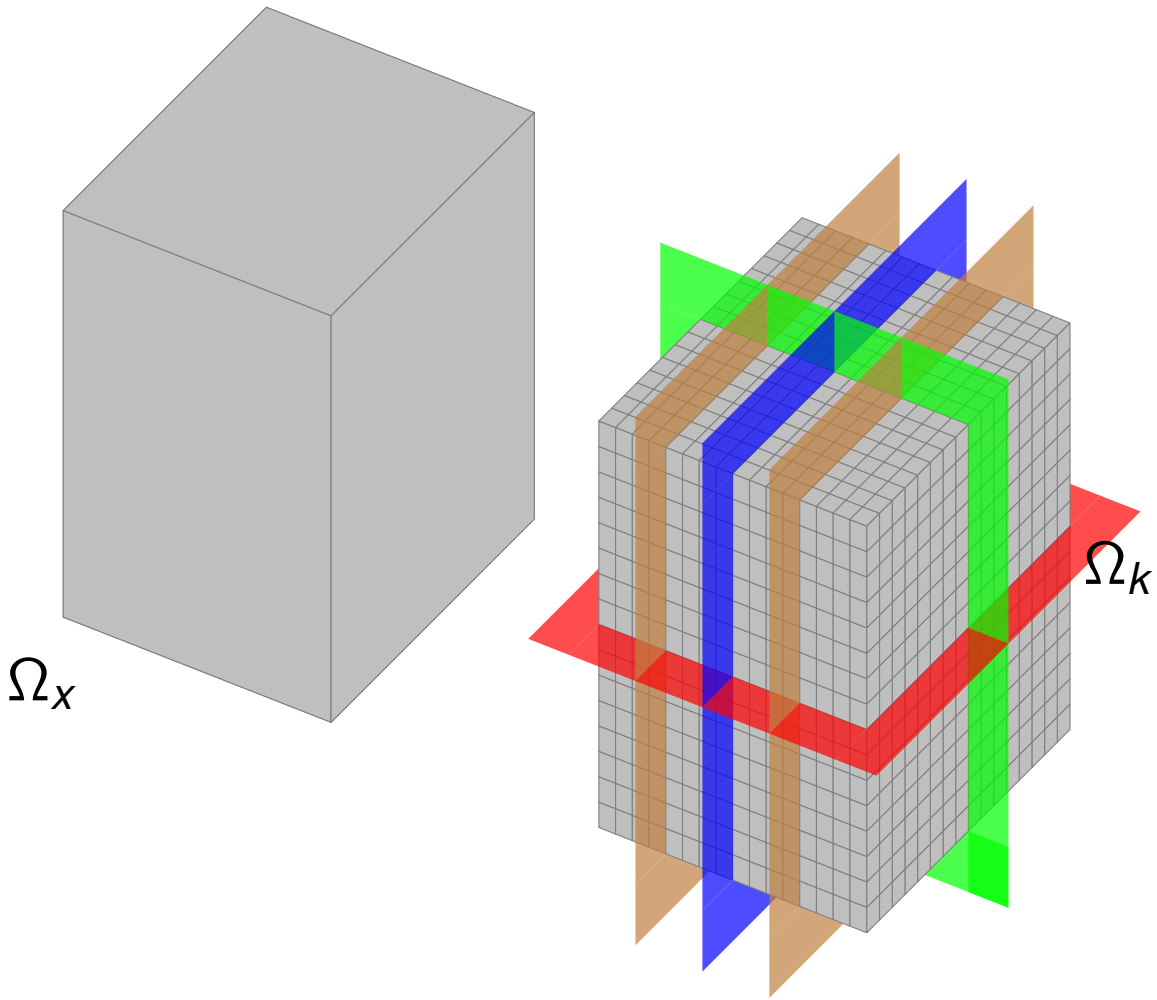
Example for $d = 3$, $N = 16$, and $p = 16$: Level 0

Each process then initializes N^d/p interactions between Ω_x and each $K \in \{K_j^0\}$ within its chunk of Ω_k.



Example for $d = 3$, $N = 16$, and $p = 16$: Level 1

Refinement of Ω_x and coarsening of Ω_k is done locally on each process since no process boundaries are crossed.



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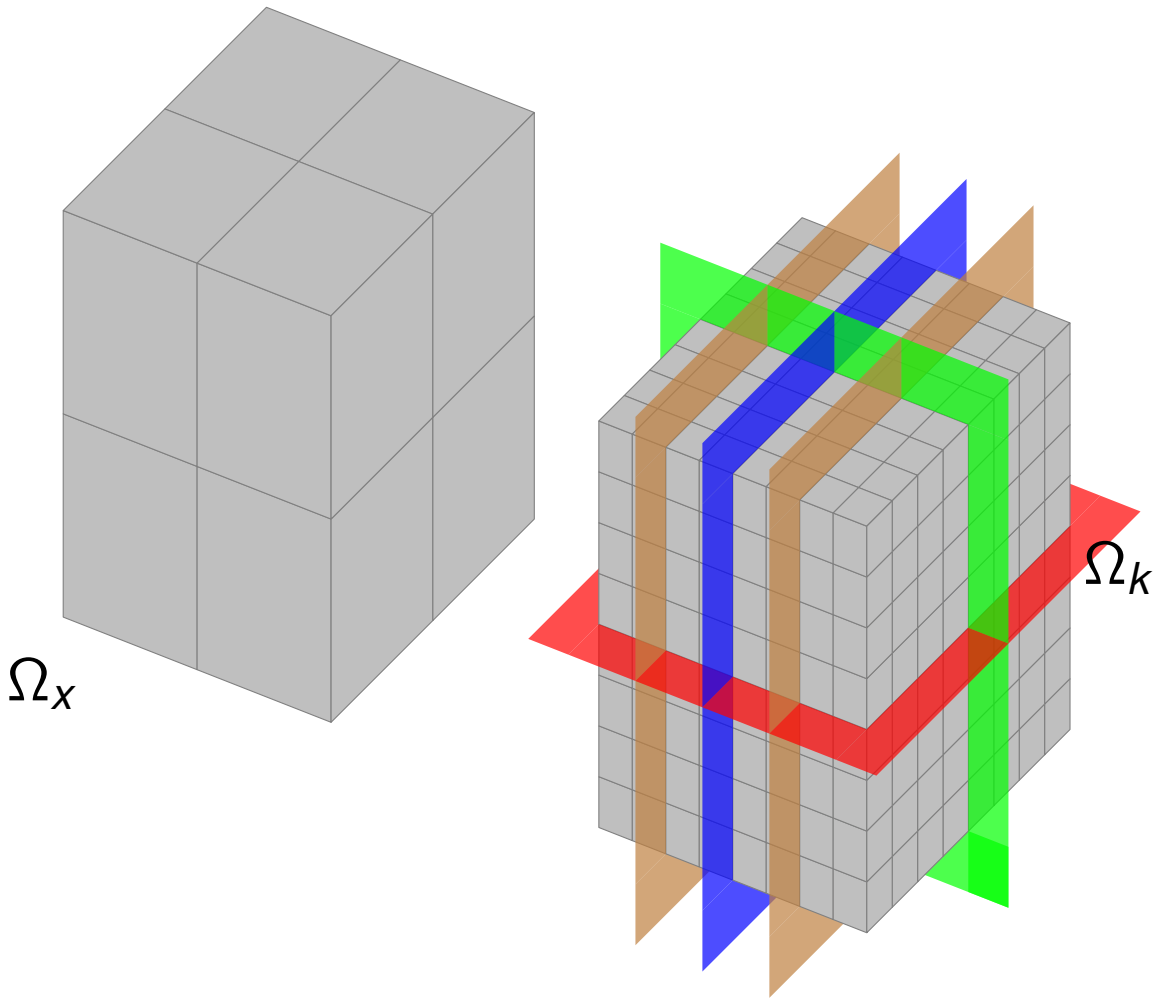
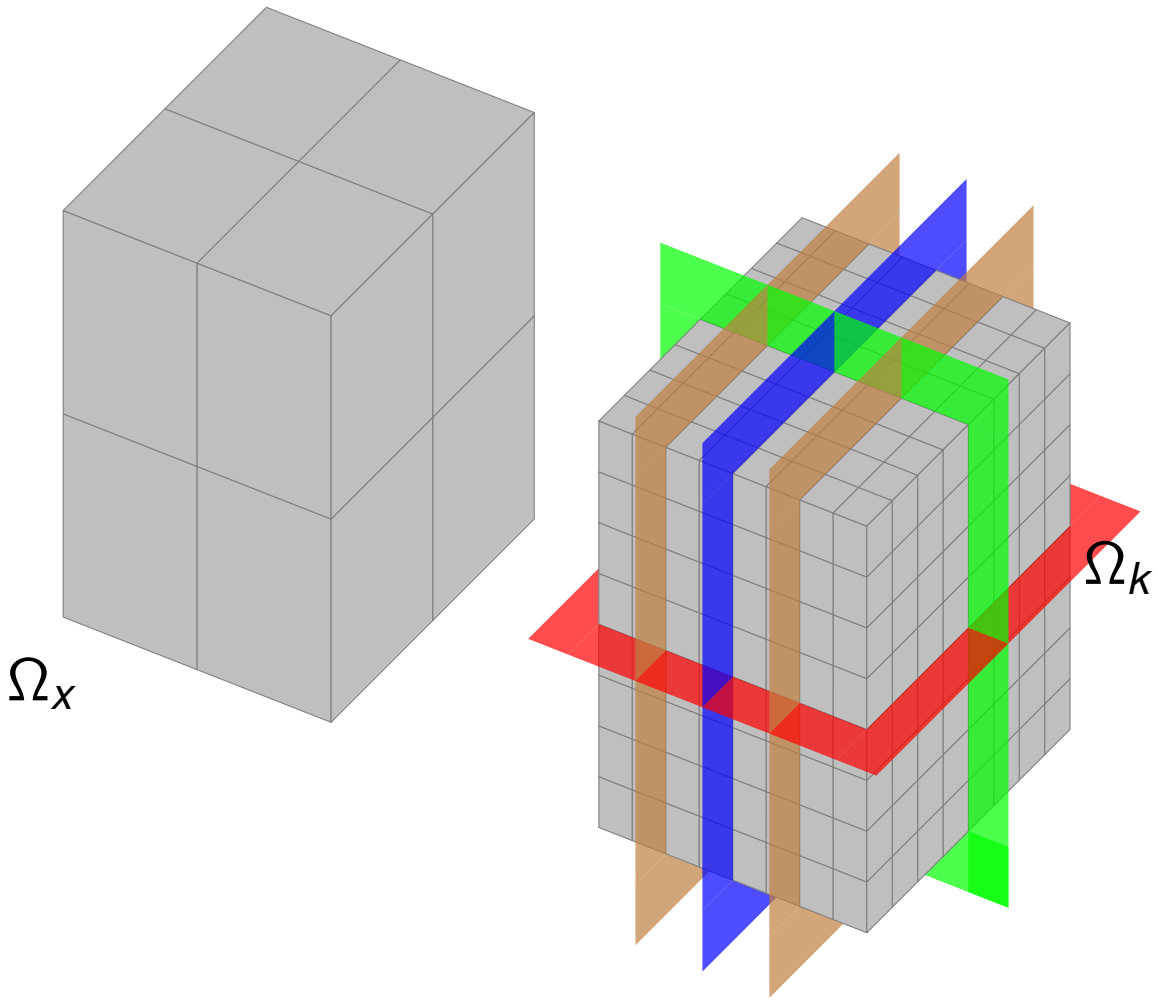


Example for $d = 3$, $N = 16$, and $p = 16$: Level 2

Refinement of Ω_x and coarsening of Ω_k is again local because no process boundaries are crossed.



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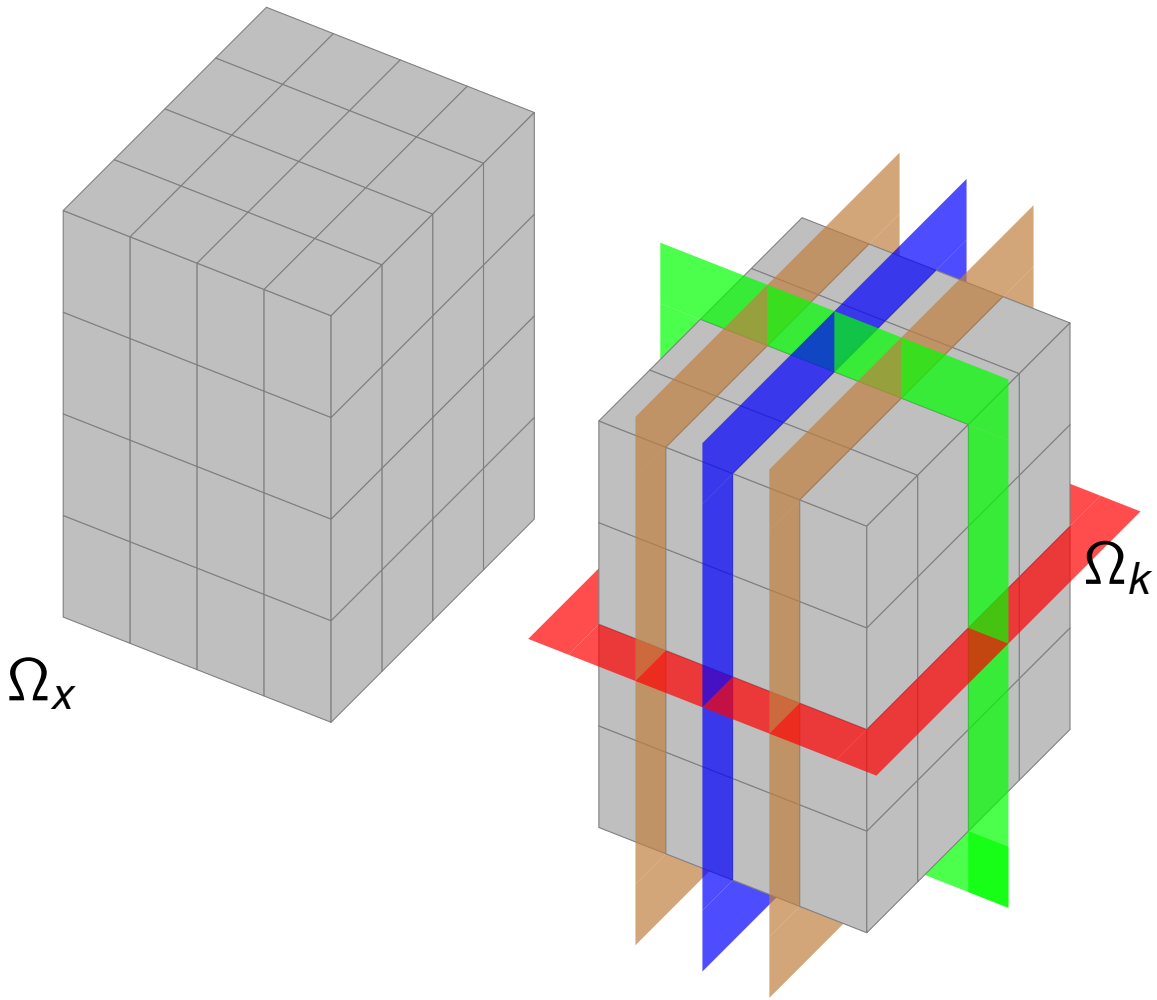
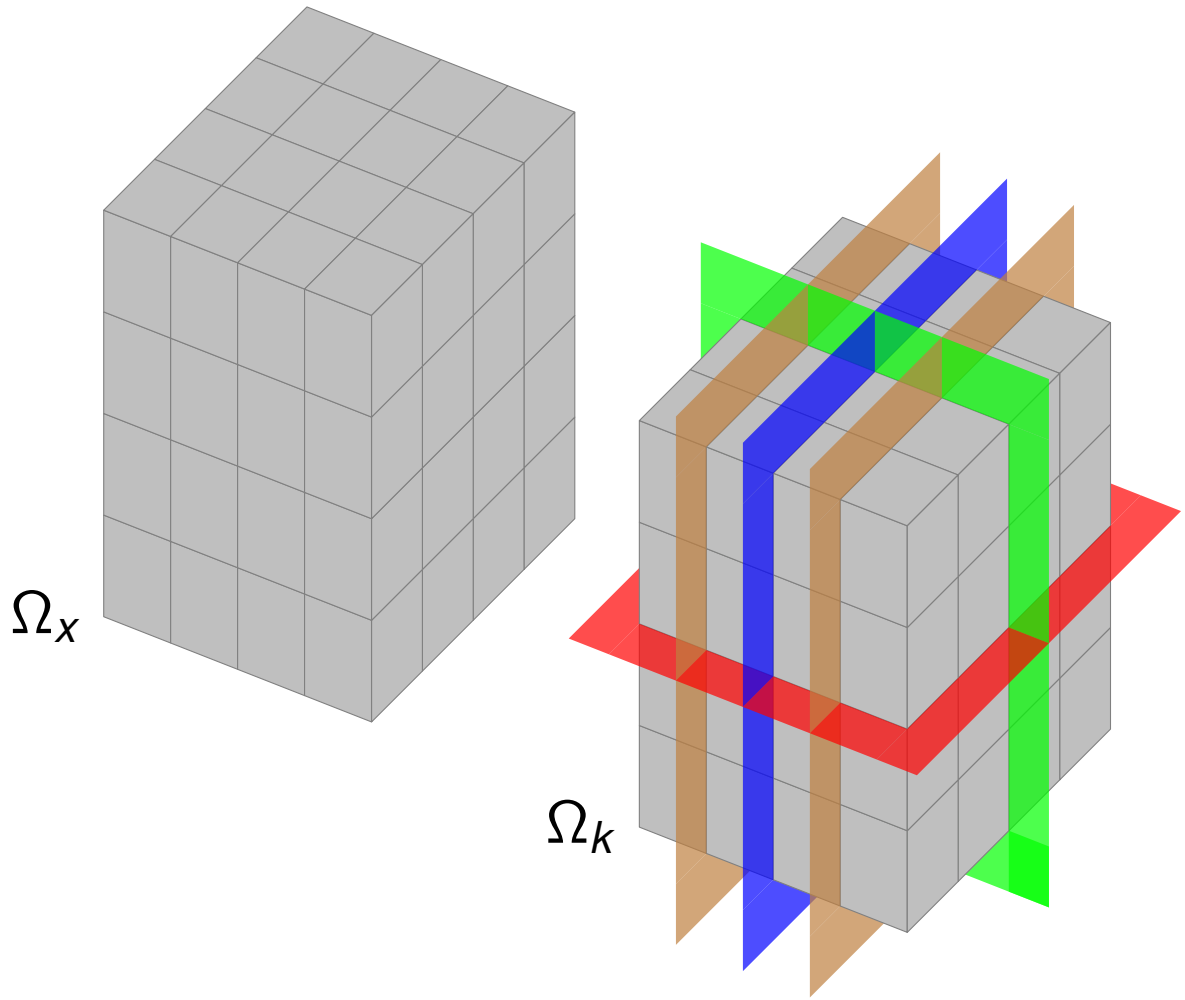


Example for $d = 3$, $N = 16$, and $p = 16$: Level 3

Coarsening Ω_k crosses the fourth set of bisections (in brown).



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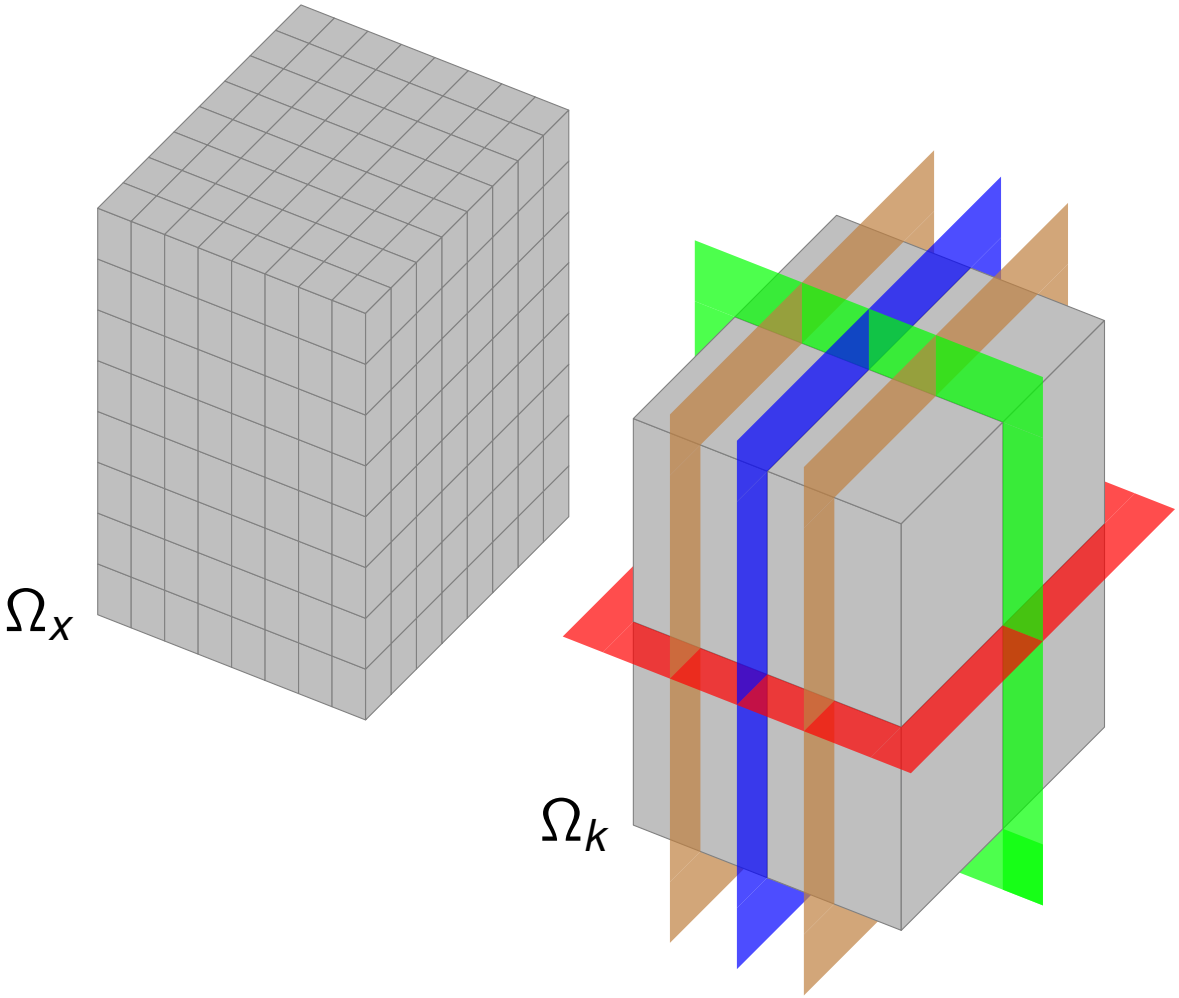
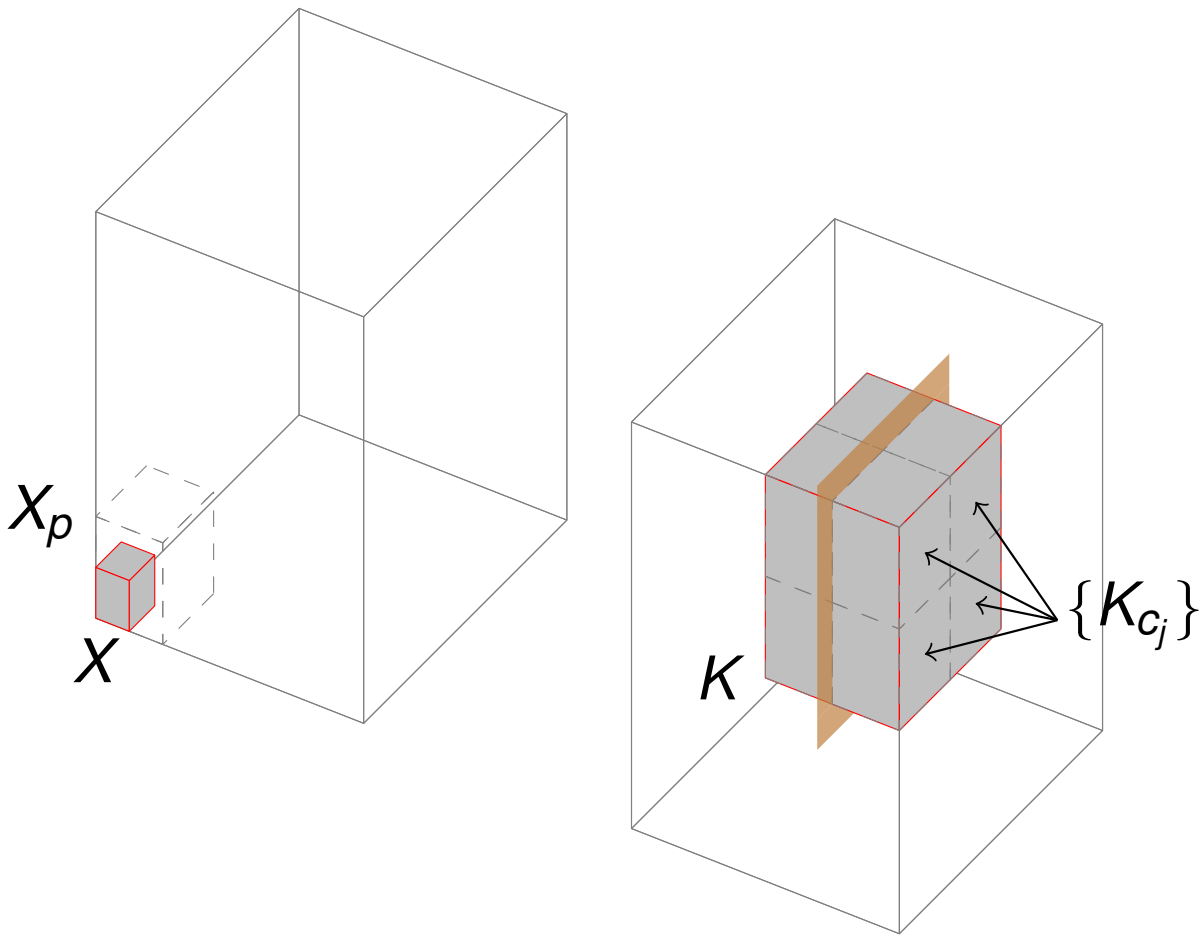


Example for $d = 3$, $N = 16$, and $p = 16$: Level 3

The $\{\delta_t^{XK}\}$ weights are a sum of the contributions from each (X_p, K_{c_j}). **MPI_Reduce** sums in teams of 2^m, m crossed parts.?



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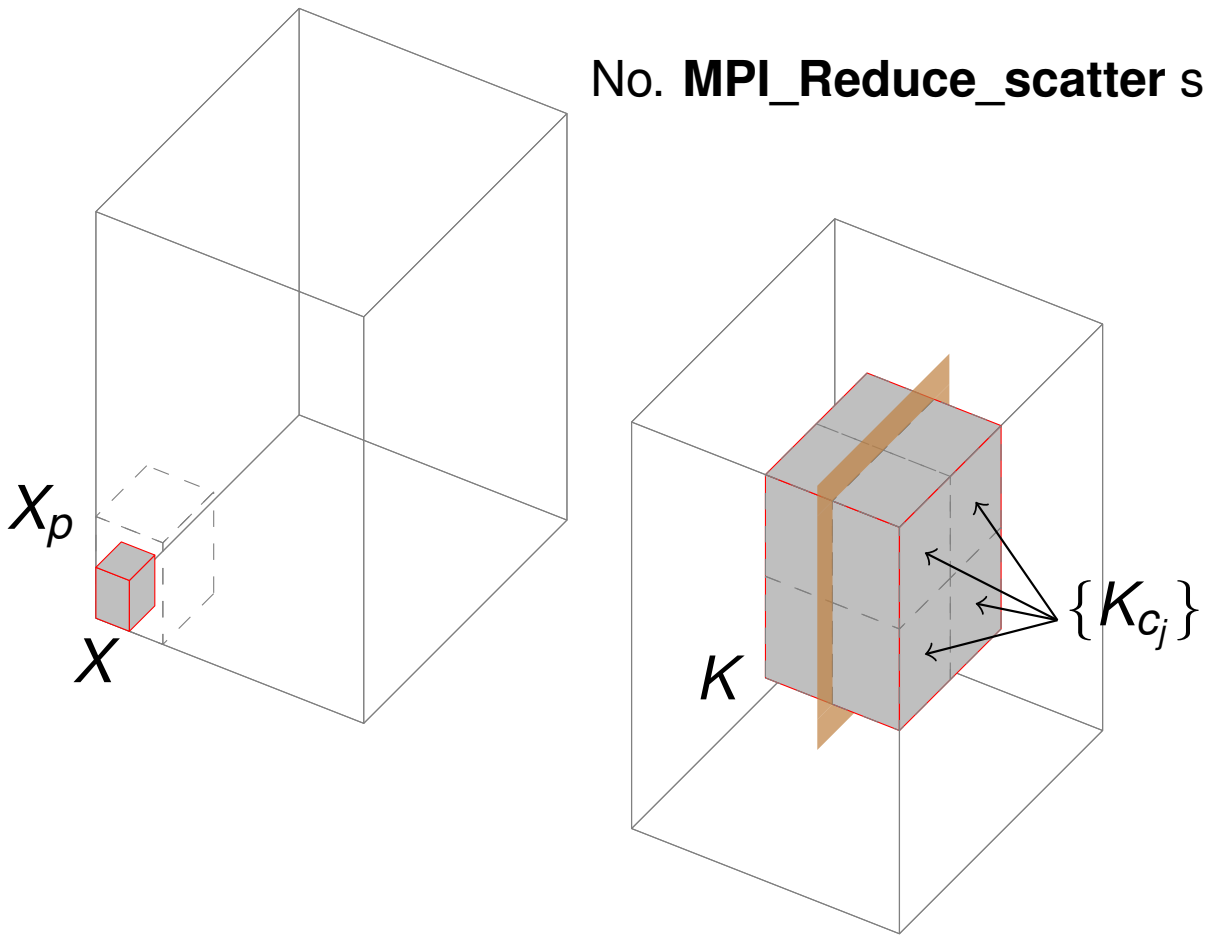
No. **MPI_Reduce_scatter** sum in teams of 2^m.



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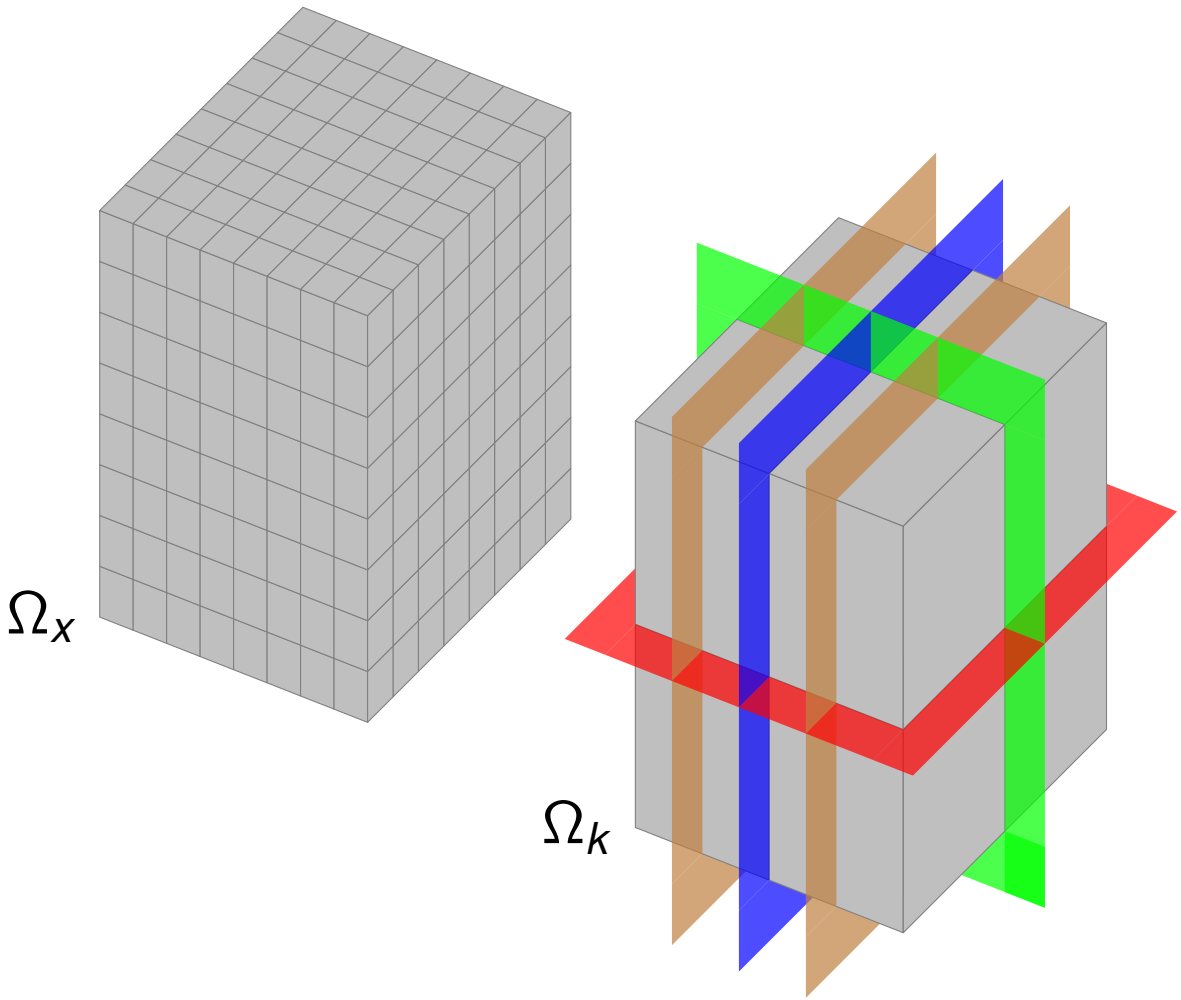
The **MPI_Reduce_scatter** is performed so that it effectively moves the crossed Ω_k partitions onto Ω_x.



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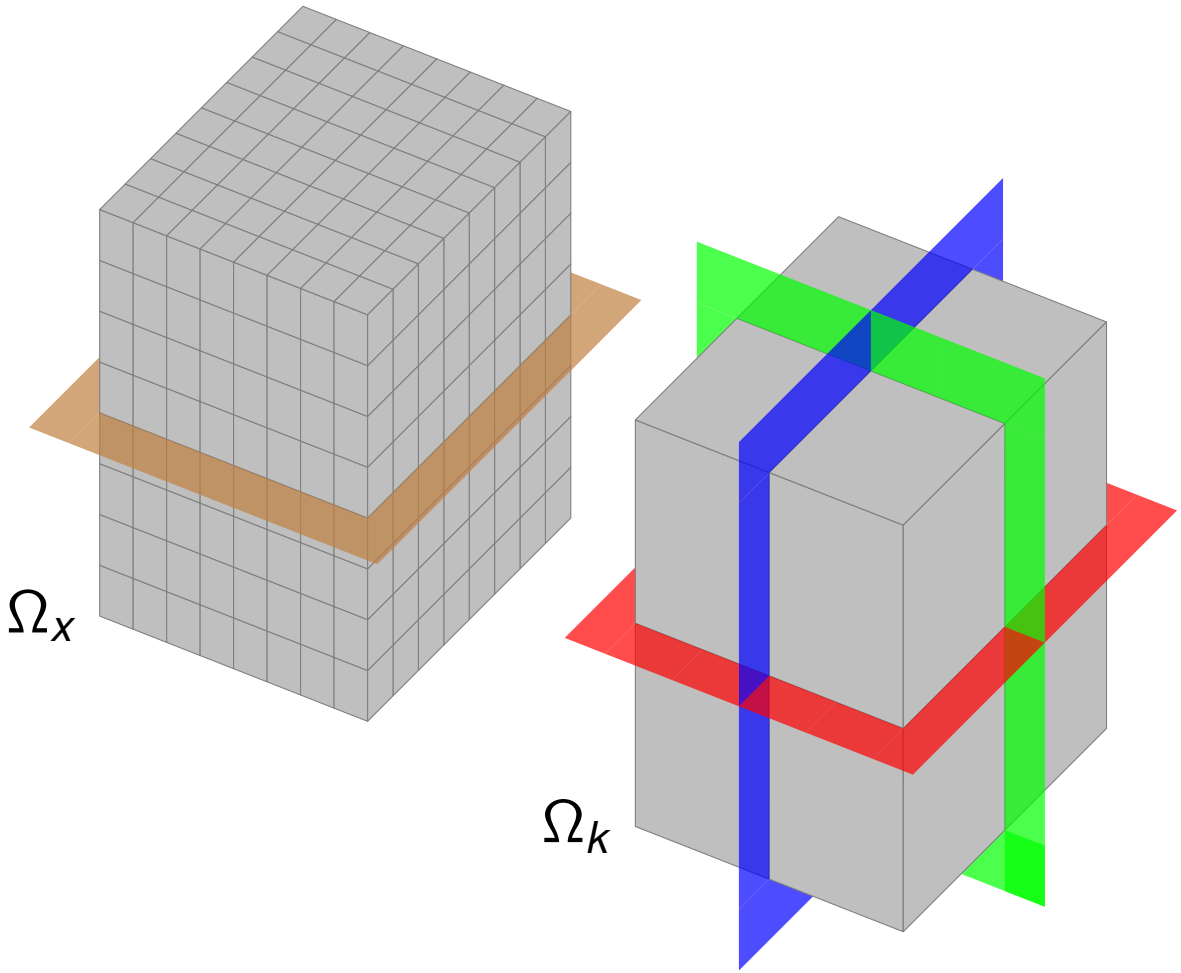
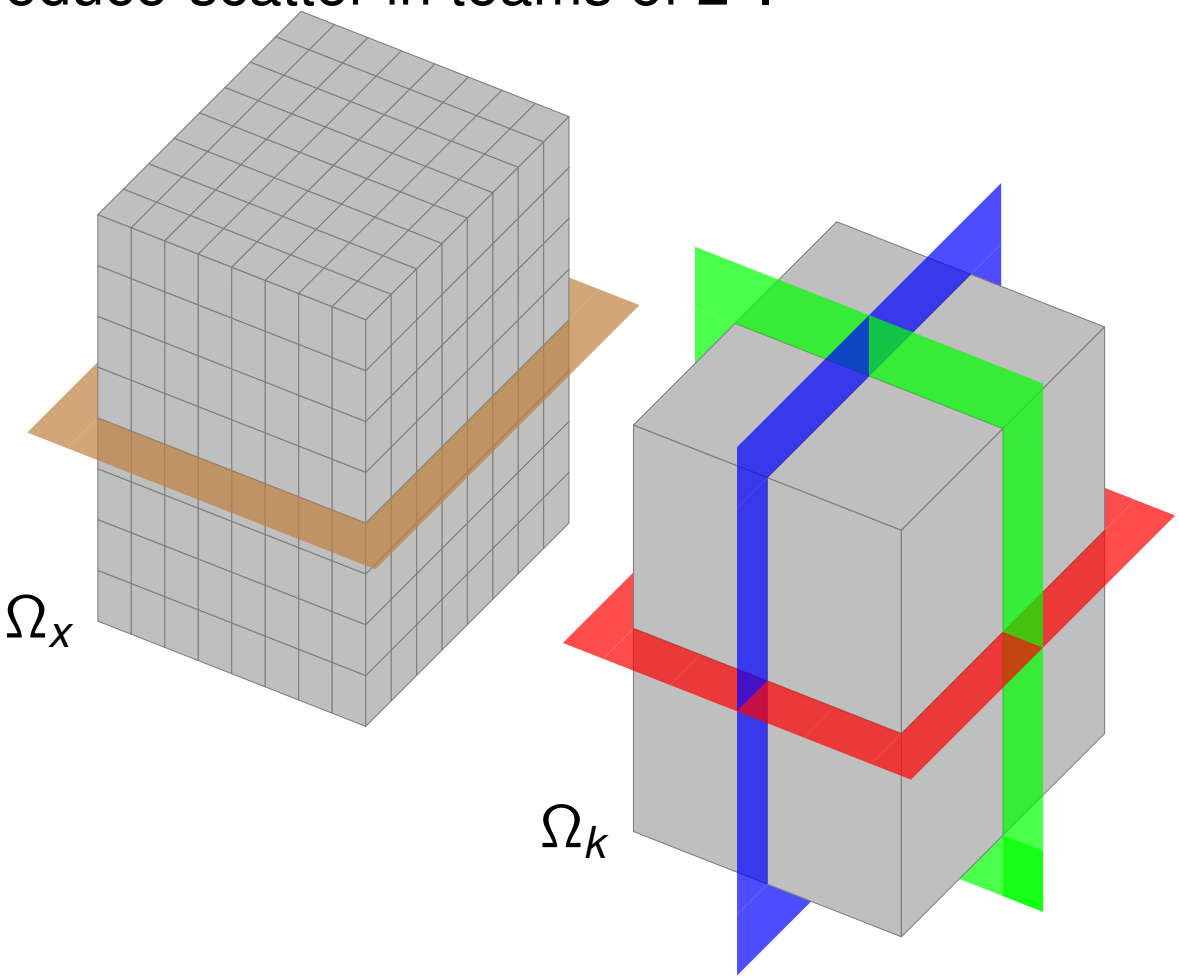


Example for $d = 3$, $N = 16$, and $p = 16$: Level 4

The last coarsening crosses 3 partition levels, so processes will reduce-scatter in teams of 2^3.



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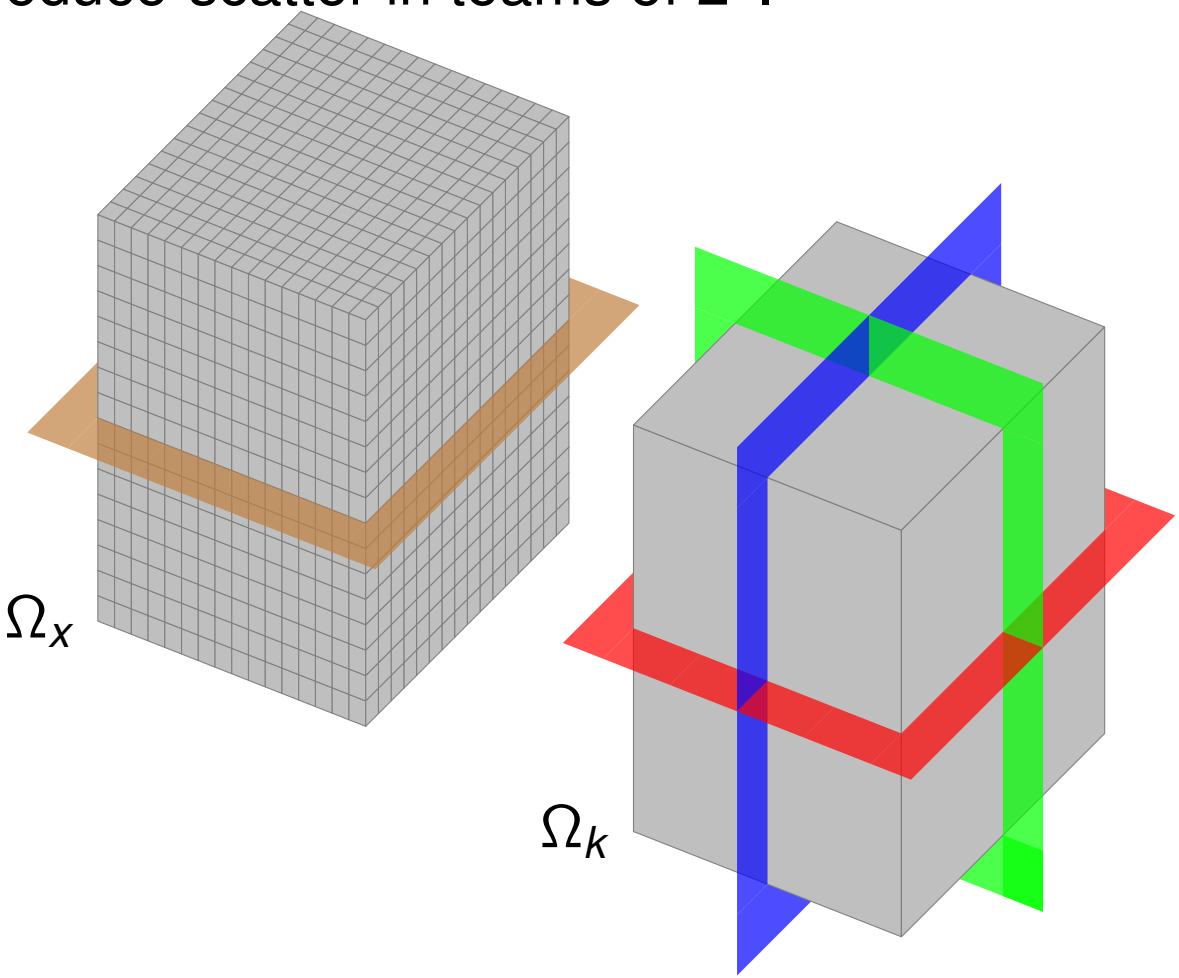

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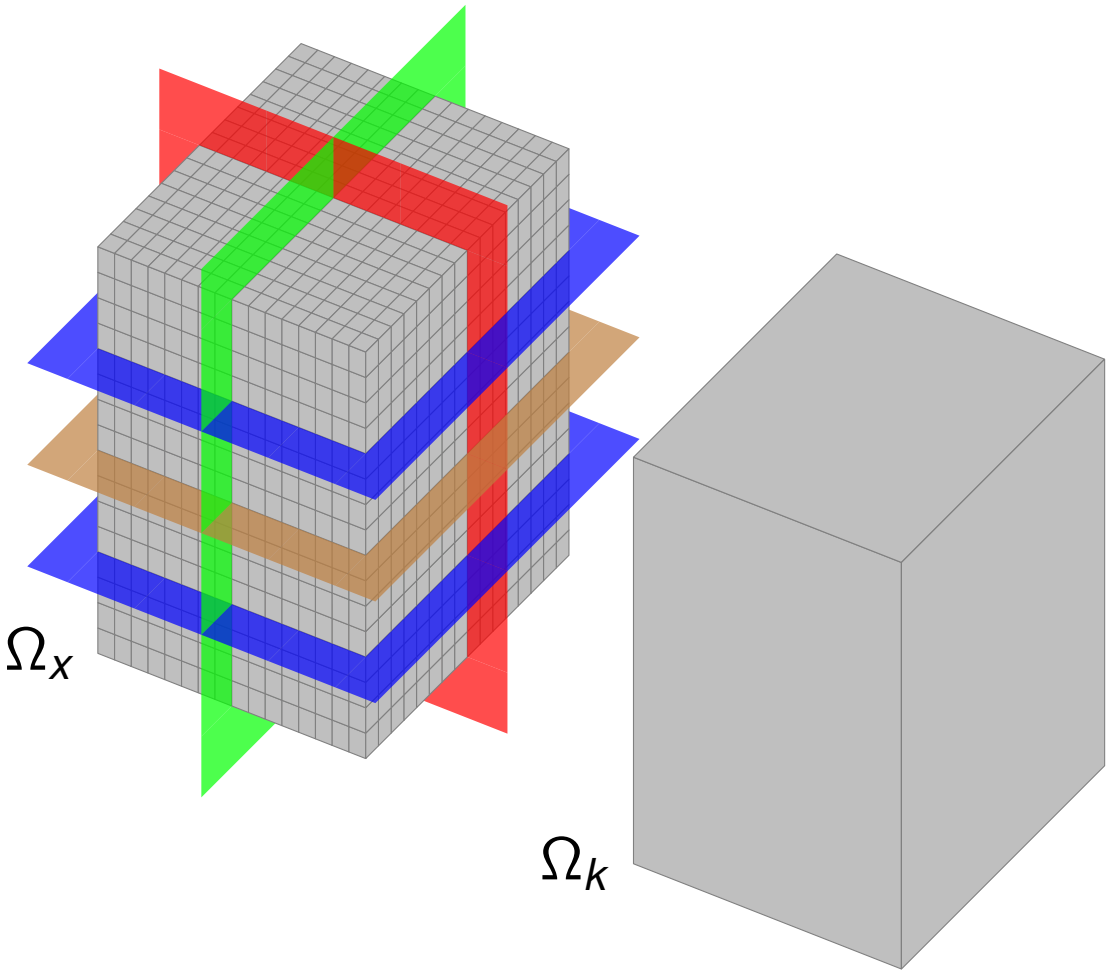
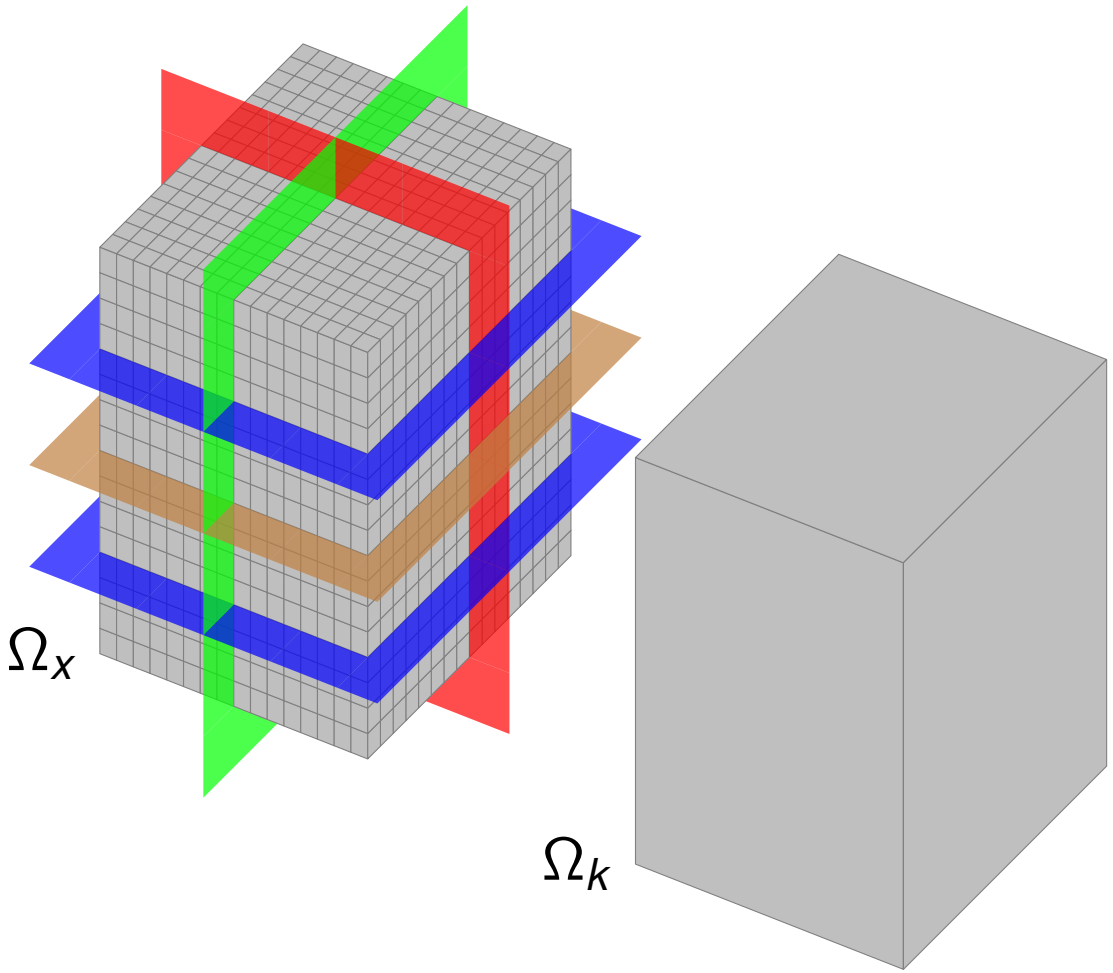


Example for $d = 3$, $N = 16$, and $p = 16$: Level 4

Done! Each process can now locally evaluate $(A_R g)(x)$ for all x within its chunk of Ω_x.



Outline

Blue Gene/P

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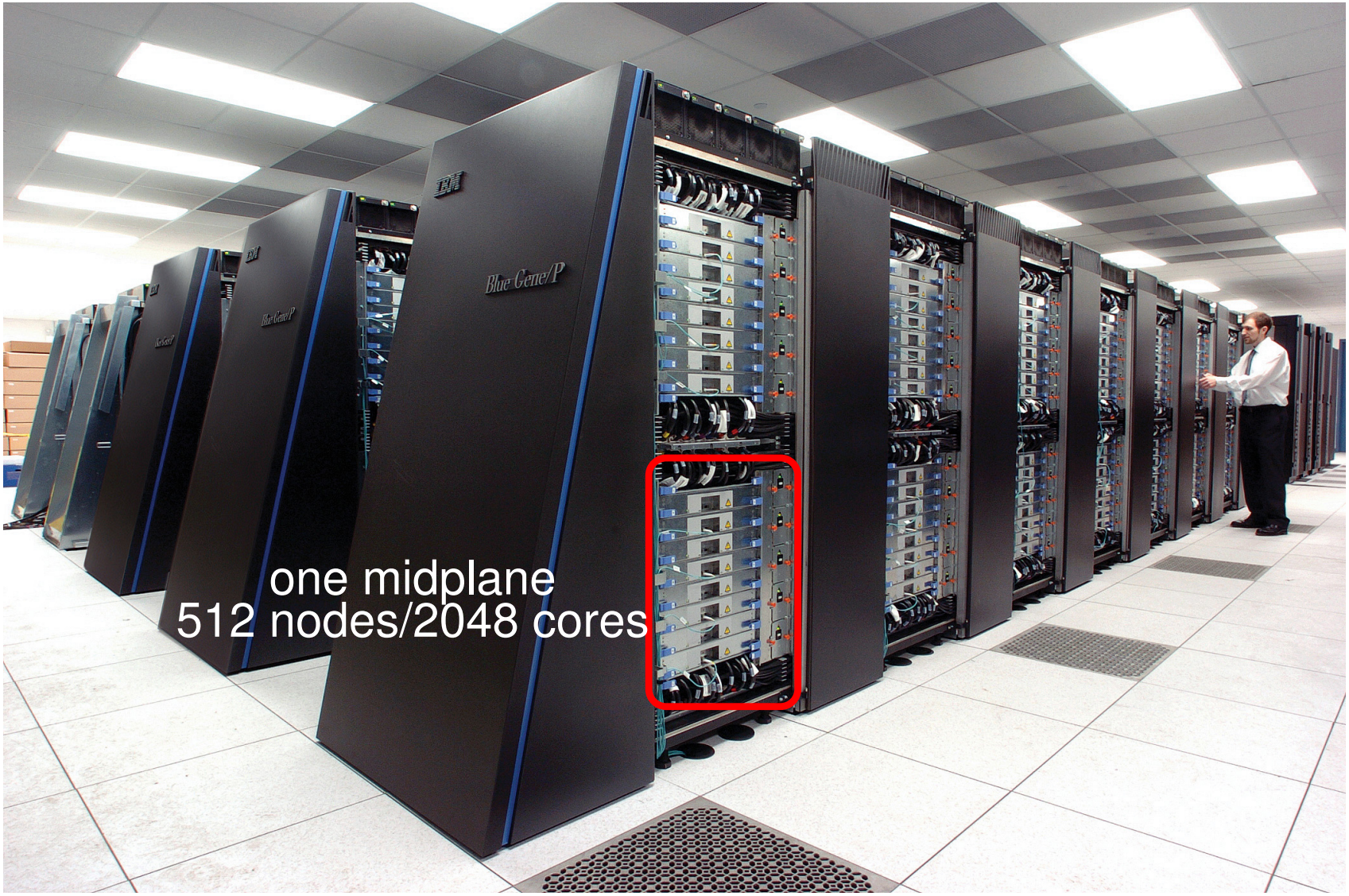
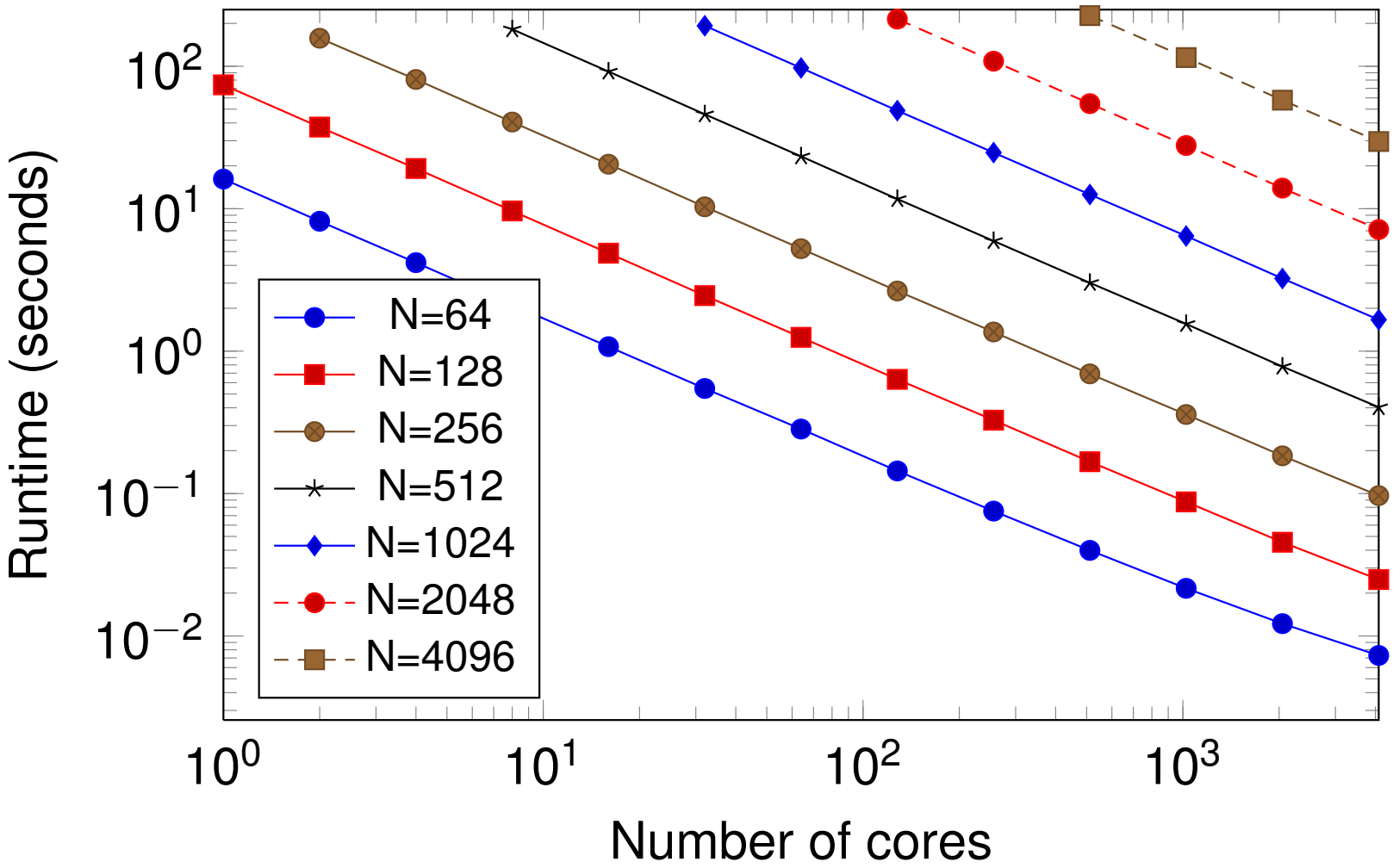
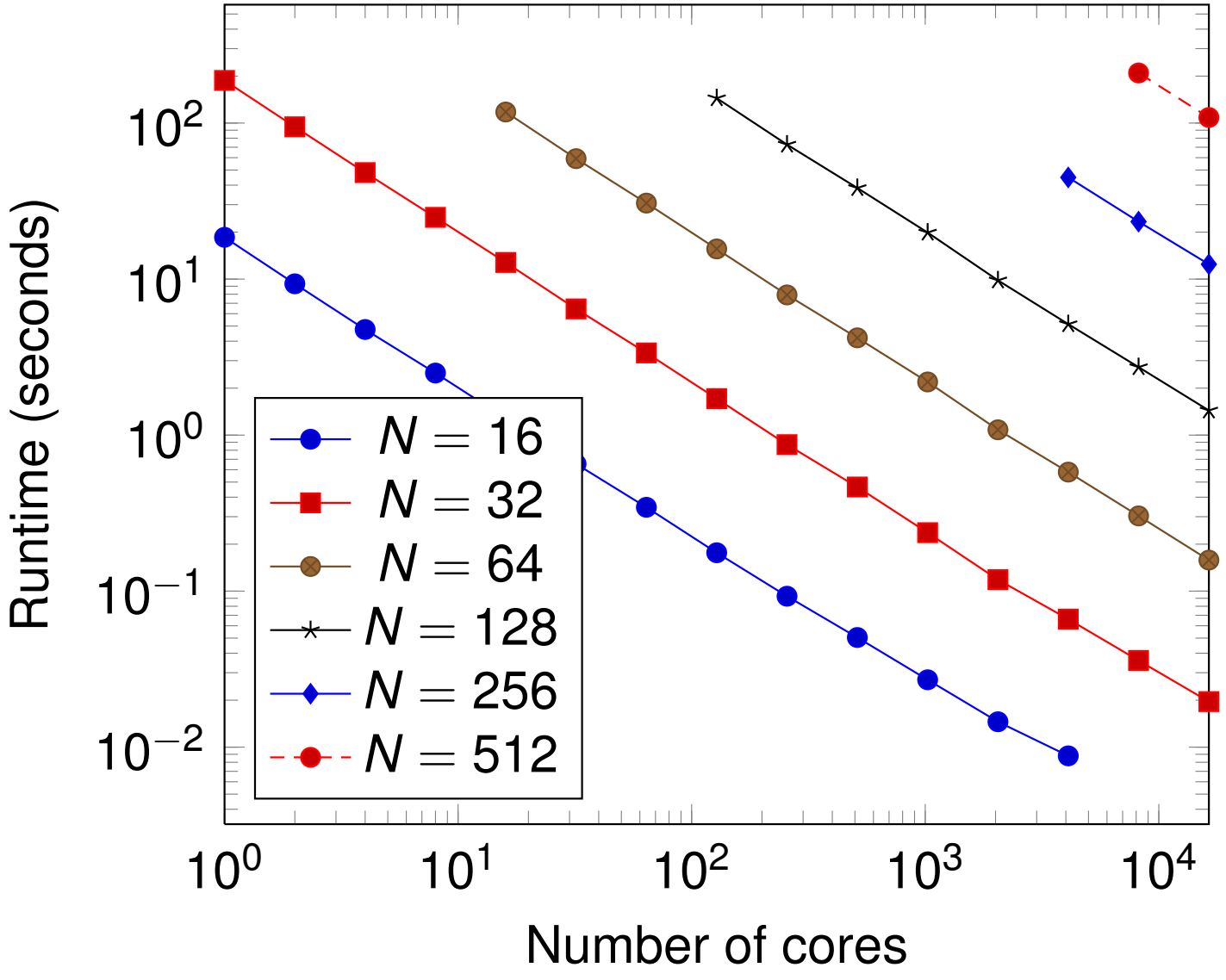


Blue Gene/P

Non-uniform 2d Generalized Radon Analogue, $q = 8$



Non-uniform 3d Fourier Transform, $q = 5$



Outline

ButterflyFIO

- Available at http://code.google.com/p/butterflyfio
- Non-uniform RFIOs and Fourier transforms already supported.
- Templated, dimension-independent, header-only C++.
- Amplitude and phase functions can simply be passed in.
- Can easily visualize potential fields with parallel VTK files.

Summary

- New parallel method for the butterfly algorithm that easily strong scales to tens of thousands of cores.
- First massively parallel fast FIO implementation.
- First massively parallel non-uniform fast Fourier transform implementation?
- Source code available at http://code.google.com/p/butterflyfio