Overview of QoS in Packet-based IP and MPLS Networks

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Agenda

- Introduction
- QoS Service Models
- DiffServ QoS Techniques
- MPLS QoS
- Summary

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What is Quality of Service?

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QoS represents the set of techniques necessary to manage network bandwidth, delay, jitter, and packet loss.

From a business perspective, it is essential to assure that the critical applications are guaranteed the network resources they need, despite varying network traffic load.

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Traffic Characterization

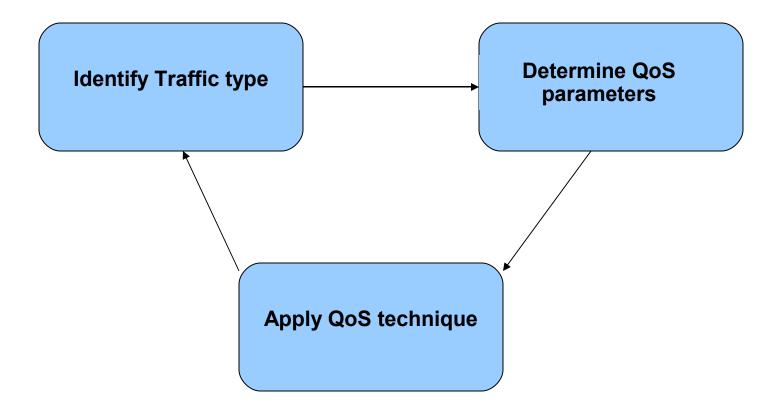
- Identify traffic sources and types
- Need for appropriate handling
 - Realtime and Non-realtime
 - Voice (Delay sensitive)
 - Video (Bandwidth intensive)
 - Data (Loss sensitive)
 - HTTP, FTP, SMTP
 - Bursty and Constant type
 - Multi-service traffic: IP, MPLS
 - Single or Multiple flows of the same type

QoS Requirements

- Traffic influencing parameters
 - Latency, Jitter, Loss
- Management of finite resources
 - Rate Control
 - Queuing and Scheduling
 - Congestion Management
 - Admission Control
 - Routing Control Traffic protection
- Service Level Agreement (SLA)
 - per-flow
 - aggregated

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QoS Triangle



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QoS Approaches

- Fine-grained approach
 - flow-based (individual flows)
- Coarse-grained approach
 - aggregated (large number of flows)
- Leads to two different QoS Models

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QoS Service Models

- Best effort (No QoS)
- Integrated services (Hard QoS)
- Differentiated services (Soft QoS)

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Best Effort Model – Traditional Internet

- "We'll do the best we can"
 But messages may be lost en route
- Traditional datagram model
- Not a traditional telephone company model
 Pay for what you want, and get exactly that

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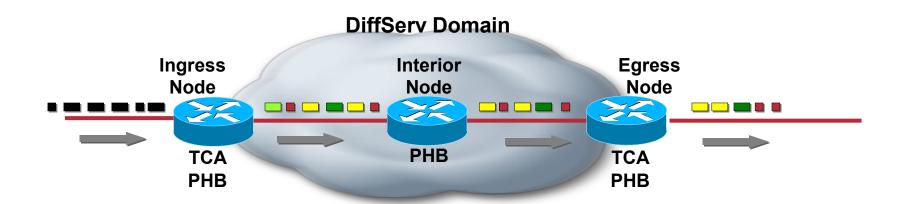
Integrated Services Model

- IntServ Architecture (RFC 1633)
- Hard QoS
- Guarantees per-flow QoS
- Strict Bandwidth Reservations
- Needs Signaling to accomplish Path Reservation
 - Resource Reservation Protocol RSVP (RFC 2205)
 - PATH/RESV messages
- Admission Control
- Must be configured on every router along the path
- Works well on small-scale
 - Has issues with scaling with large number of flows
 - Requires devices to retain state information

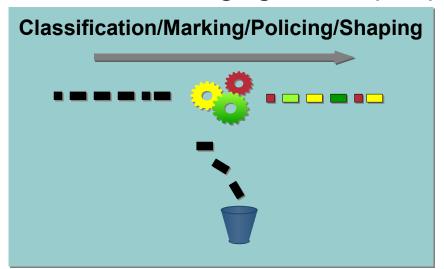
Differentiated Services Model

- DiffServ Architecture RFC 2475
- Scales well with large flows through aggregation
- Creates a means for traffic conditioning (TC)
- Defines per-hop behavior (PHB)
- Edge nodes perform TC
 - Allows core routers to do more important processing tasks
- Tough to predict end-to-end behavior
 - Especially with multiple DiffServ Domains
 - DiffServ implementation versus Capacity planning

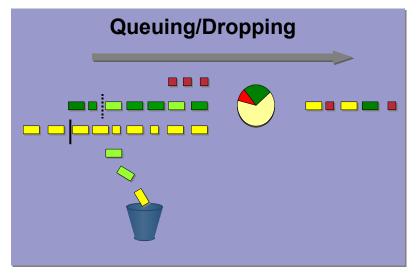
Differentiated Services Architecture



Traffic Conditioning Agreement (TCA)



Per-Hop Behavior (PHB)



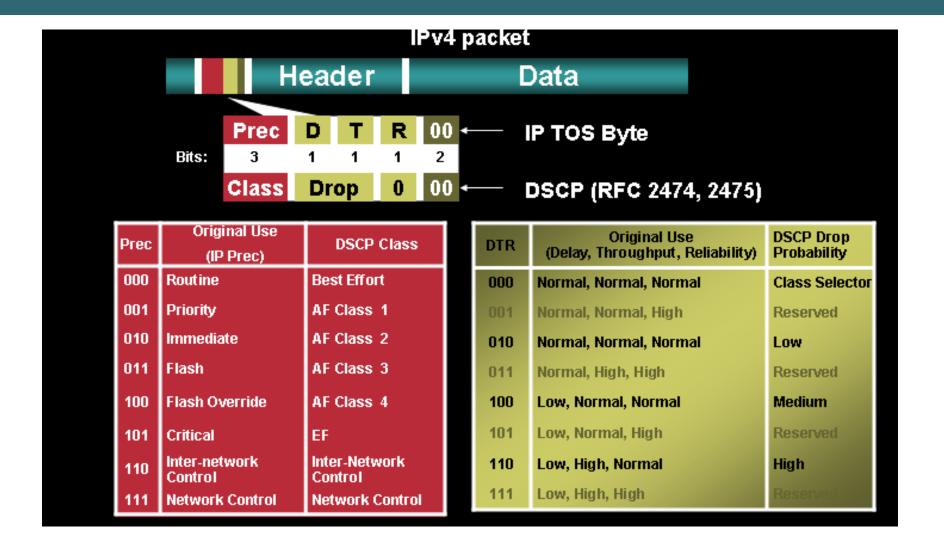
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IETF DiffServ Model

- Re-define TOS byte in IP header to Differentiated Services Code Point (DSCP)
- Uses 6 bits to categorize traffic into "Behavior Aggregates"
- Defines a number of "Per Hop Behaviors" applied to links
- Two-Ingredient Recipe:
 - Condition the Traffic at the Edges
 - Invoke the PHBs in the Core

IP TOS vs IP DSCP



Diffserv Class Selector

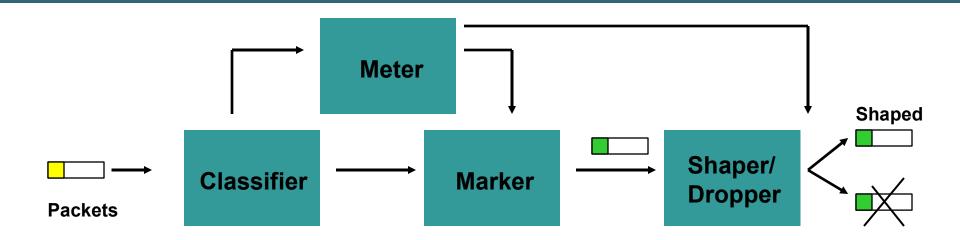
Class Selector provides support for IP Prece using DSCP terminology

	Туре	Class Selector Code Point
Prec 0	Routine	000 <mark>000</mark> (0)
Prec 1	Priority	001 <mark>000</mark> (8)
Prec 2	Immediate	010 <mark>000</mark> (16)
Prec 3	Flash	011000 (24)
Prec 4	Override	100 <mark>000</mark> (32)
Prec 5	Critical	101 <mark>000</mark> (40)
Prec 6	Inter-net	110000 (48)
Prec 7	Net-Control	111 <mark>000</mark> (56)

"match ip dscp 24"
is the same as
"match ip precedence 3"

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DiffServ Traffic Conditioner



- Classifier: Selects a packet in a traffic stream based on the content of some portion of the packet header
- Meter: Checks compliance to traffic parameters (eg Token Bucket) and passes result to the marker and shaper/dropper to trigger a particular action for in/out of profile packets
- Marker: Writes/rewrites DSCP
- Shaper: Delays some packets to be compliant with a profile
- Dropper: Discards some or all of the packets in a traffic stream in order to bring the stream into compliance with a traffic profile

Classification and Marking

Classification

- Identification based on field(s) in a packet
- Flow identification parameters
 - Src/Dest. Address, Source/Dest. Port, Protocol
- IP Precedence / DSCP based

Marking

- Marking/Coloring packets to indicate class
- Application marked or node configured
 - IP Precedence or DSCP
 - MPLS EXP
 - Other instances (FR-DE and ATM-CLP)

Traffic Metering

- Traffic Rate Management in network boundary nodes
- Traffic Metering measures traffic
 - Does not alter traffic characteristics
 - Reports compliance results to Shaper or Dropper
- Uses Token Bucket Scheme to measure traffic
 - Mean or Committed Information Rate
 - Conformed Burst size
 - Extended Burst size

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Policing and Shaping

Police

Sends conforming traffic and allows bursts

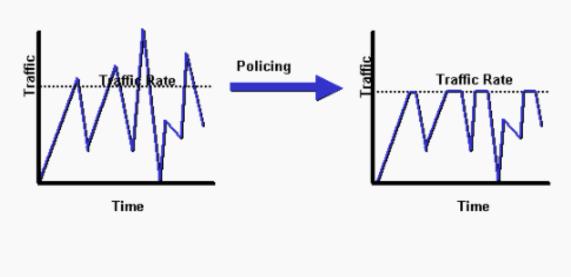
Drops non-conforming traffic (due to lack of tokens)

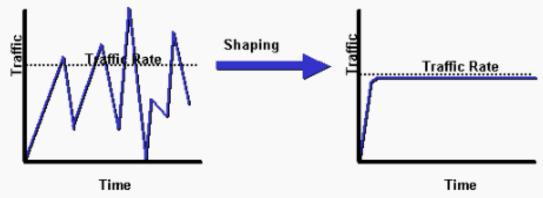
Provision for Packet re-marking

Shaping

- Smoothes traffic but increases overall latency
- Buffers packets when tokens are exhausted

Policing and Shaping

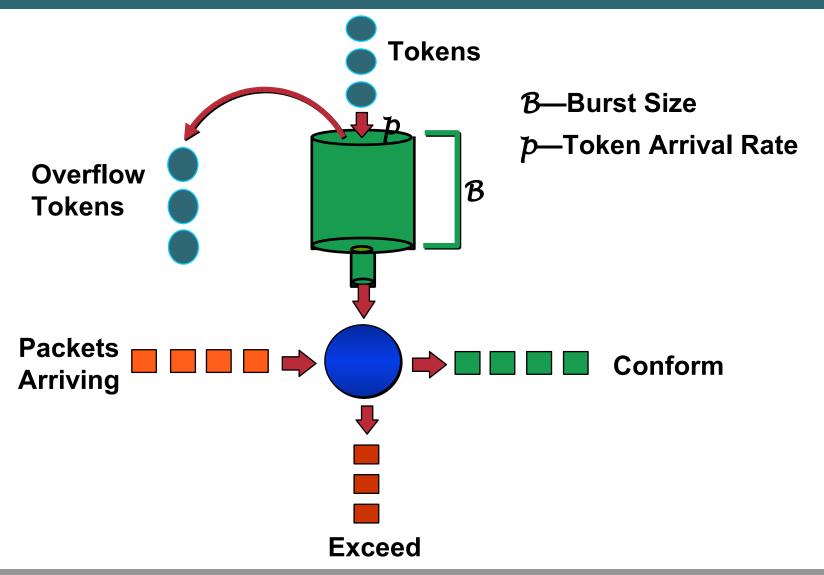




Policing

- Uses the token bucket scheme
- Tokens added to the bucket at the committed rate
- Depth of the bucket determines the burst size
- Packets arriving with sufficient tokens in the bucket are said to conform
- Packets arriving with insufficient tokens in the bucket are said to exceed

Token Bucket in Policing

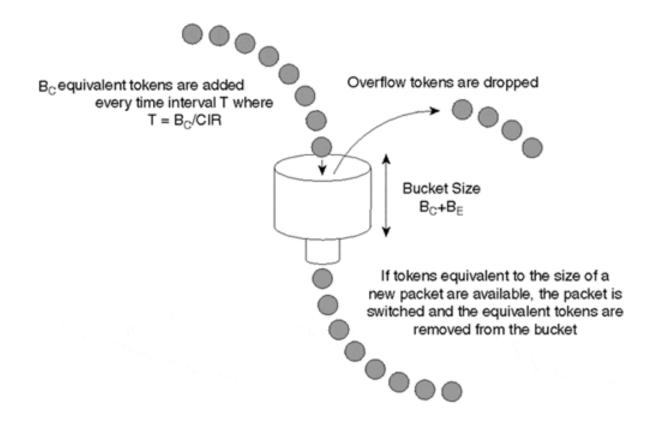


Shaping

- Uses the token bucket scheme
- Smoothes bursty traffic to meet CIR through buffering
- Queued Packets transmitted as tokens are available

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Token Bucket in Traffic Shaping



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Per-Hop Behavior (PHB)

- PHB relates to resource allocation for a flow
- Resource allocation is typically Bandwidth
- Queuing / Scheduling mechanisms:
 - FIFO / WFQ / MWRR / MDRR
- PHB also involves determining a packet drop policy
- Congestion avoidance schemes primary technique RED / WRED

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Queuing/Scheduling

- Scheduling mechanisms guarantee BW for flows
- More bandwidth guarantee means dequeue more from one queue or set of queues.
- De-queue depends on weights allocated to queues

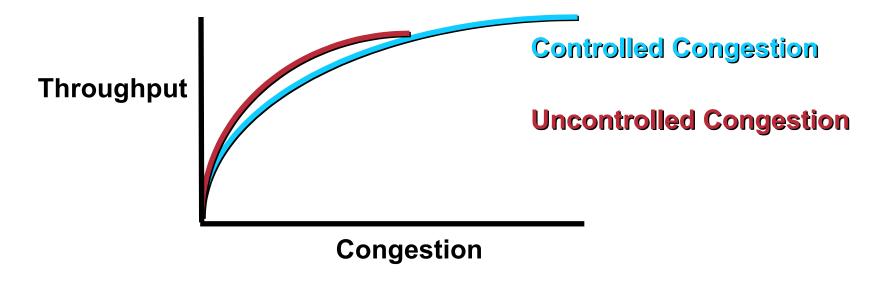
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Congestion Avoidance/Management

- When there is congestion what should we do?
 Tail drop i.e. Packets dropped due to Max Queue Length
 Drop selectively but based on IP Prec / DSCP bit
- Congestion control mechanisms for TCP traffic
 - Adaptive
 - Dominant transport protocol

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The Problem of Congestion



Uncontrolled, congestion will seriously degrade performance

The system buffers fill up

Packets are dropped, resulting in retransmissions

This causes more packet loss and increased latency

The problem builds on itself

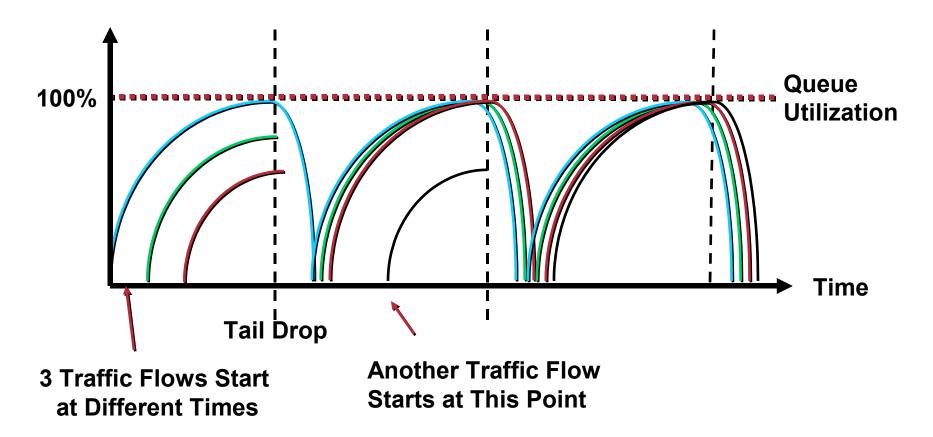
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TCP traffic and Congestion

- Congestion window based on slow-start
 Sender / Receive negotiation
- Packet loss indicator of congestion
 - Congestion window re-sizing
 - Source throttles traffic

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Global Synchronization



Global synchronization is many connections going through TCP Slow-Start mode at the same time

Random Early Detect (RED)

- A congestion avoidance algorithm
- Designed to work with a transport protocol like TCP
- Minimize packet delay jitter by controlling average queue size
- Uses Packet drop probability and Avg. Queue size
- Avoids global synchronization of many connections

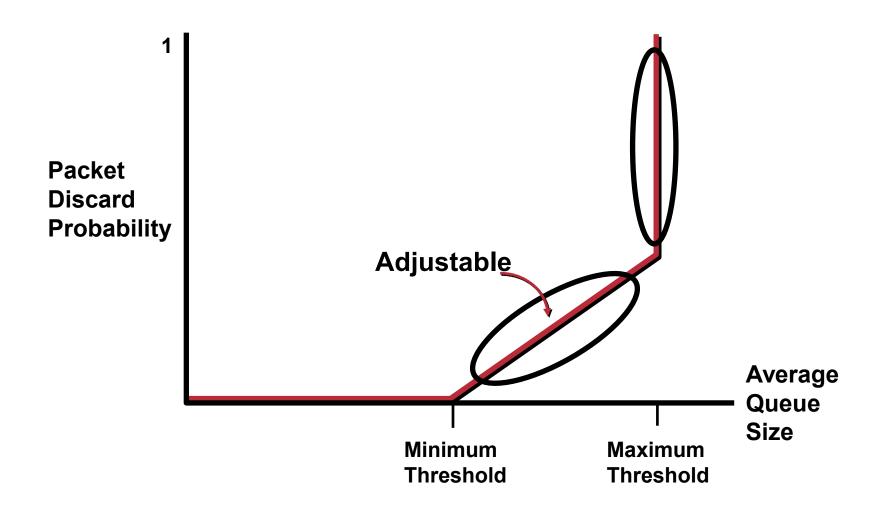
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RED—Packet-Drop Probability

- Packets are dropped sufficiently frequently to control the average queue size
- The probability that a packet is dropped from a connection is proportional to the amount of packets sent by the connection

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RED



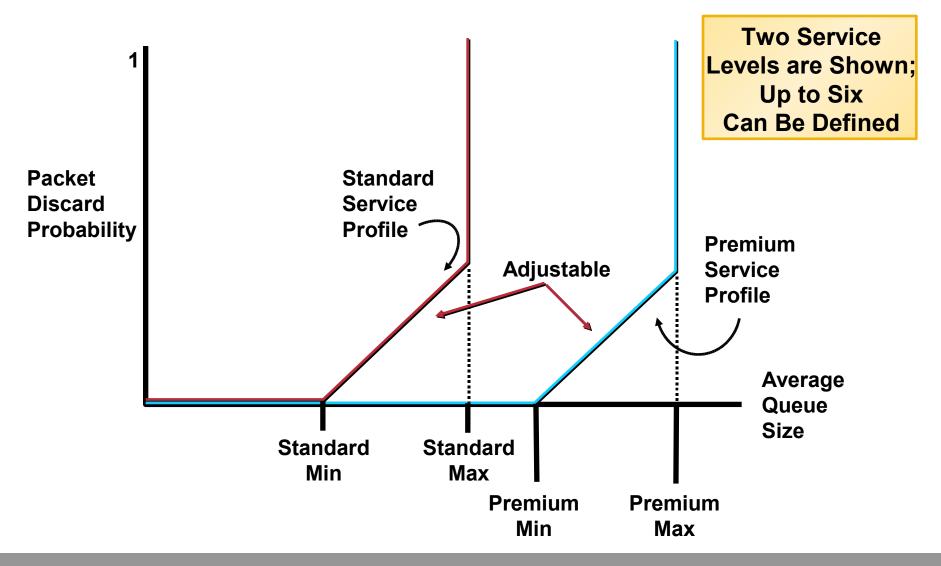
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Weighted RED (WRED)

- WRED combines RED with IP Precedence or DSCP to implement multiple service classes
- Each service class has a defined min and max thresholds, and drop rates

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WRED Service Profile Example



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When Should WRED be Used?

 Where the bulk of your traffic is TCP as oppose to UDP

Only TCP will react to a packet drop; UDP will not

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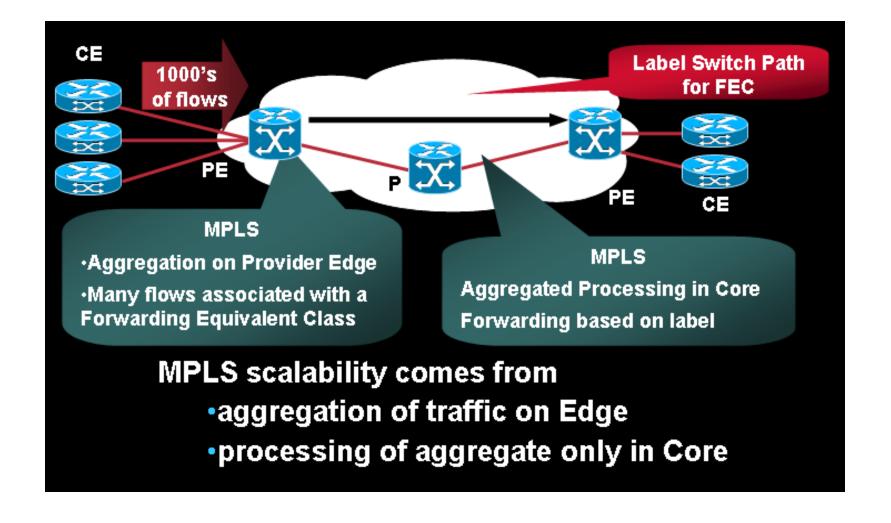
MPLS Diffserv

MPLS DiffServ Architecture

- MPLS does NOT define new QoS architectures
- MPLS QoS uses Differentiated Services (DiffServ) architecture defined for IP QoS (RFC 2475)
- MPLS DiffServ is defined in RFC3270

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DiffServ Scalability via Aggregation



What's Unchanged in MPLS DiffServ

When Compared to IP DiffServ

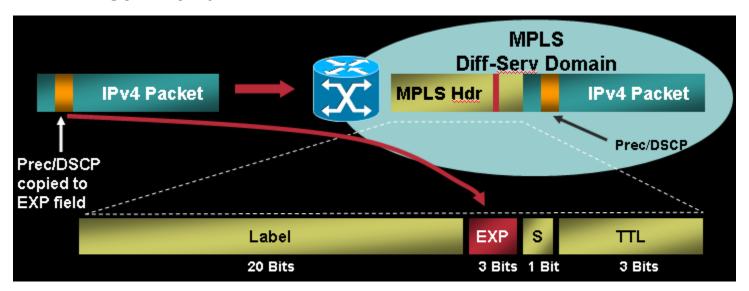
Functional components (TCA/PHB) and where they are used
 Classification, marking, policing, and shaping at network
 boundaries

Buffer management and packet scheduling mechanisms used to implement PHB

- PHB definitions
 - EF: low delay/jitter/loss
 - AF: low loss
 - BE: No guarantees (best effort)

What's new in MPLS DiffServ?

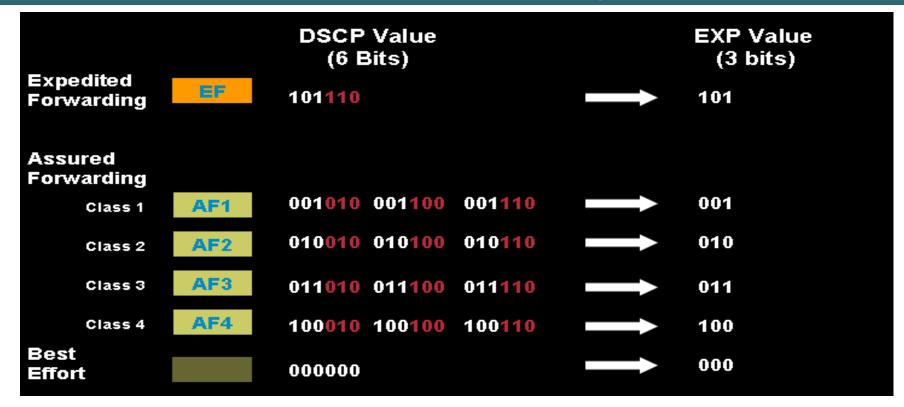
IP DiffServ Domain



- Prec/DSCP field is not directly visible to MPLS Label Switch Routers (they forward based on MPLS Header and EXP field)
- Information on DiffServ must be made visible to LSR in MPLS Header using EXP field / Label.
- How do we map DSCP into EXP? Interaction between them.

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DSCP to EXP Mapping



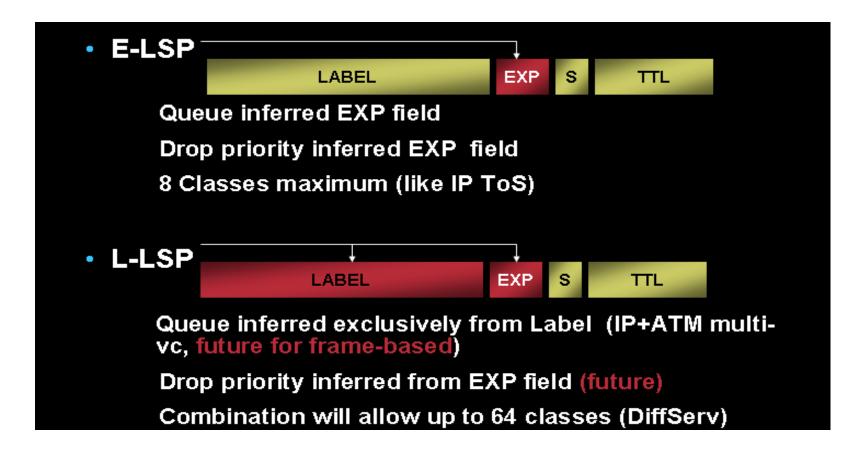
RFC3270 does not recommend specific EXP values for DS PHBs (EF/AF/CS)

MPLS DiffServ – RFC 3270

- Problem: IP DSCP = 6 bits while MPLS EXP = 3bits
- Solution: where 8 or less PHBs are used, those can be mapped into EXP field→use "E-LSPs with preconfigured mapping"
- Solution: where more than 8 PHBs are used in core, those need to be mapped in both "label and EXP" → "L-LSPs" are needed

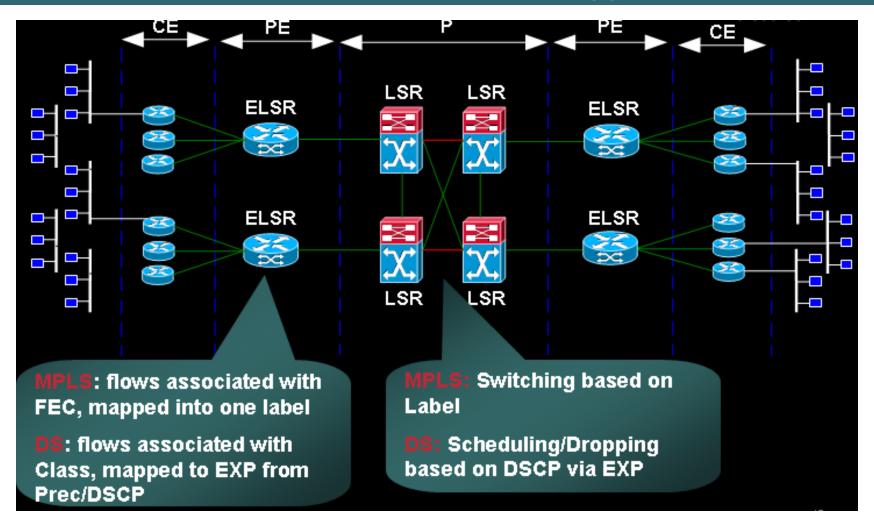
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Types of Label Switched Paths



Both E-LSP and L-LSP can use LDP or RSVP for label distribution

MPLS DiffServ Topology



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MPLS DiffServ Tunneling Modes

- Based on RFC 3270
- Modes

Uniform

Short-Pipe

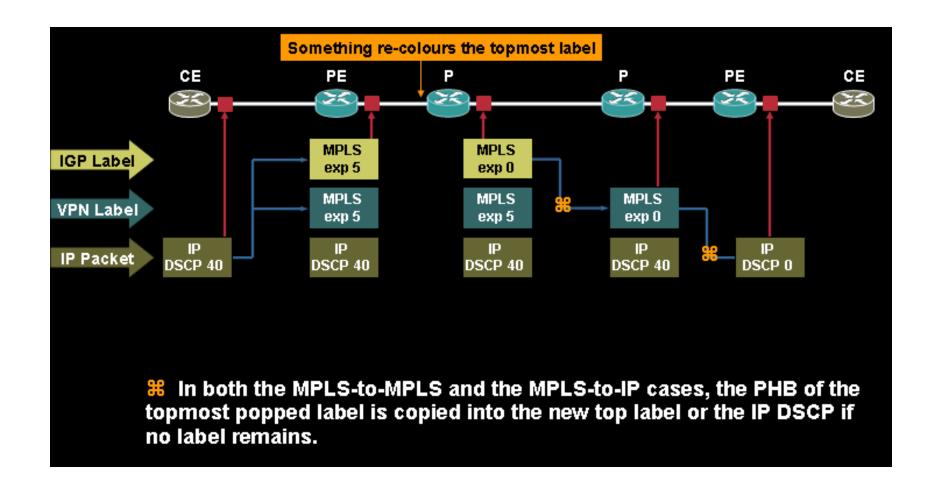
Pipe

Uniform Mode

- Assume the entire admin domain of a Service Provider is under a single DiffServ domain
- Then, it is likely a requirement to keep the colouring information uniform (keep it when going from IP to IP, IP to MPLS, MPLS to MPLS, MPLS to IP).

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Uniform Mode



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Short-Pipe Mode

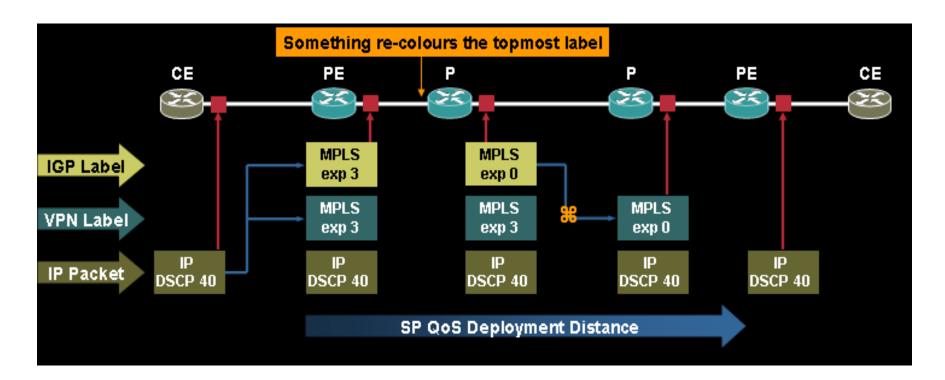
- Assume an ISP network implementing a DiffServ Policy
- Assume its customer network implementing another policy
- Requirement:

Transparency: the customer wants to preserve its DSCP intact

Uniformity: within the IP/MPLS backbone, the SP wants to have a uniform diffserv domain

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Short-Pipe Mode



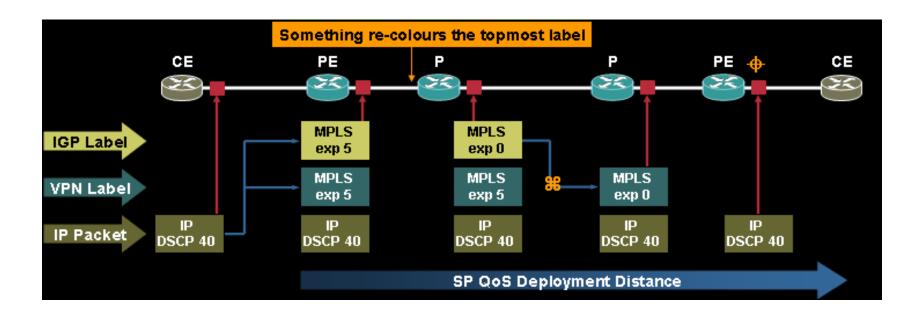
- Hereby The PHB of the topmost popped label is copied into the new top label
- Note that policy applied on the egress interface of the egress PE is based on the DSCP of the customer, hence the 'short-pipe' naming.

Pipe Mode

- Exactly the same case as Short-Pipe
- However, the SP wants to drive the outbound classification for WFQ/WRED on the egress interface from a PE to a CE based on its DiffServ policy (EXP)

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Pipe Mode



- Here is a superior of the topmost popped label is copied into the new top label
- Classification is based on mpls-exp field (EXP=0) of the topmost received MPLS frame

MPLS TE and Diffserv

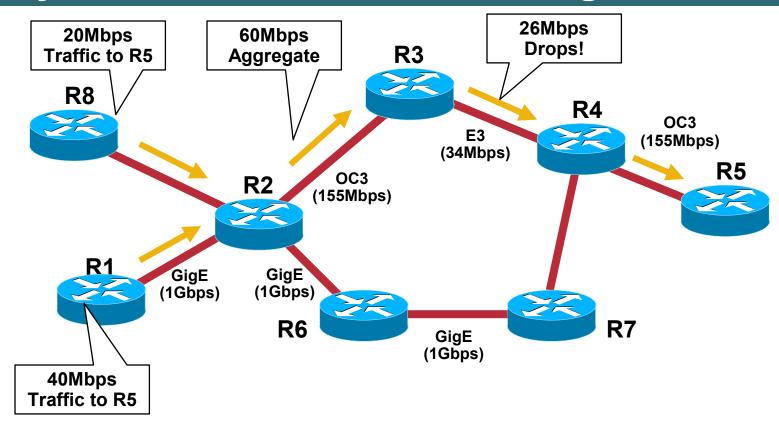
BW Optimization and Congestion Mgmt. in Parallel

TE + DiffServ

- Spread Traffic around with more flexibility than the IGP Offers
- Reserve per-class bandwidth, sort of
- Manage Unfairness During Temporary Congestion

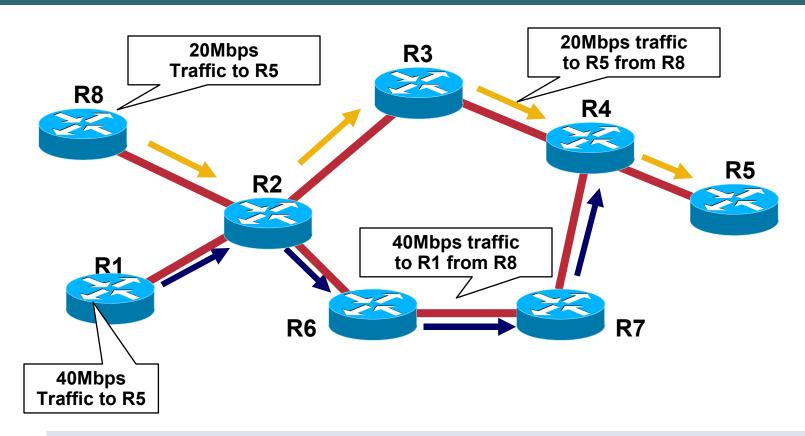
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Why TE: Shortest Path and Congestion



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The TE Solution

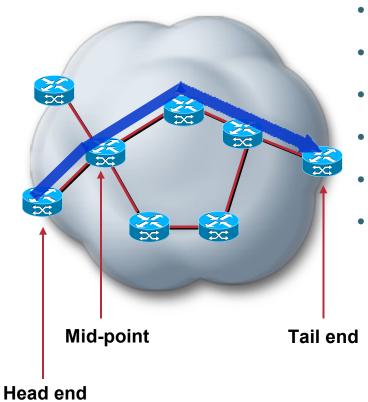


- MPLS Labels can be used to engineer explicit paths
- Tunnels are UNI-DIRECTIONAL

Normal path: R8 \rightarrow R2 \rightarrow R3 \rightarrow R4 \rightarrow R5

Tunnel path: R1 \rightarrow R2 \rightarrow R6 \rightarrow R7 \rightarrow R4

How MPLS TE Works



- Explicit routing
- Constrained-based routing
- Admission control
- Protection capabilities
- RSVP-TE to establish LSPs
- ISIS and OSPF extensions to advertise link attributes

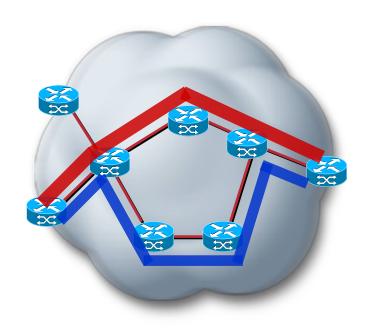
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DiffServ-Aware TE (DS-TE)

- Regular TE allows for one reservable bandwidth amount per link
- DS-TE allows for more than one reservable bandwidth amount per link
- Brings per-class dimension to TE
- Basic idea: connect PHB class bandwidth to DS-TE bandwidth sub-pool

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DiffServ-Aware TE

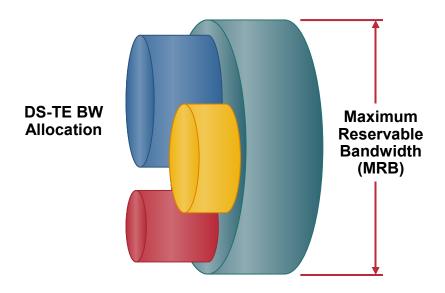


- Per-class constrained-based routing
- Per-class admission control



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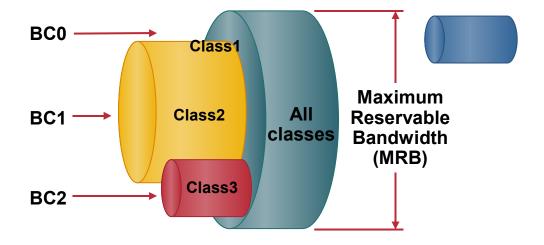
DiffServ-Aware TE



- Link BW distributed in pools or BW Constraints (BC)
- Up to 8 BW pools
- Different BW pool models

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DS-TE BW Pools – Maximum Allocation Model (MAM)



BC0: 20% Best Effort

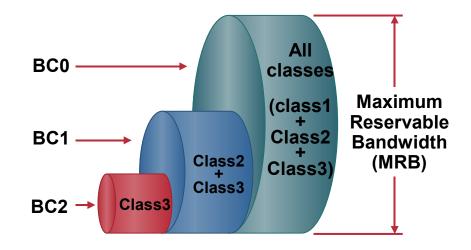
BC1: 50% Premium

BC2: 30% Voice

DS-TE BW Pools – Russian Dolls Model (RDM)

- BW pool applies to one or more classes
- Global BW pool (BC0) equals MRB
- BC0..BCn used for computing unreserved BW for class n





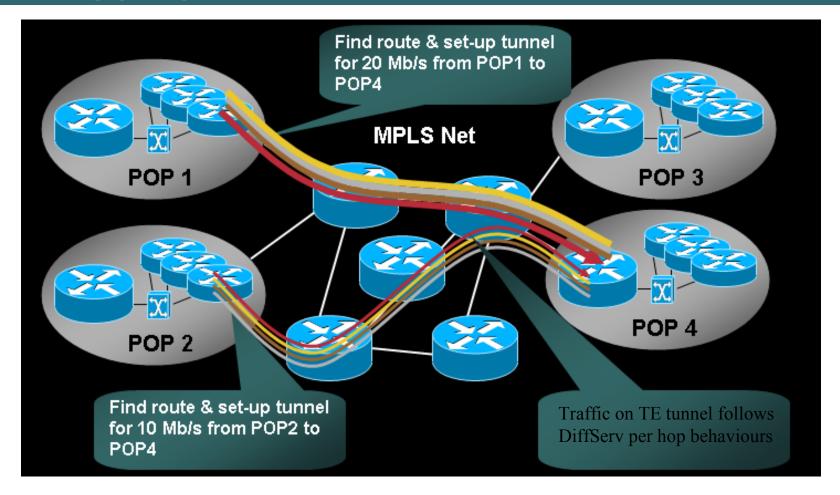
BC0: MRB Best Effort + Premium + Voice

BC1: 50% Premium + Voice

BC2: 30% Voice

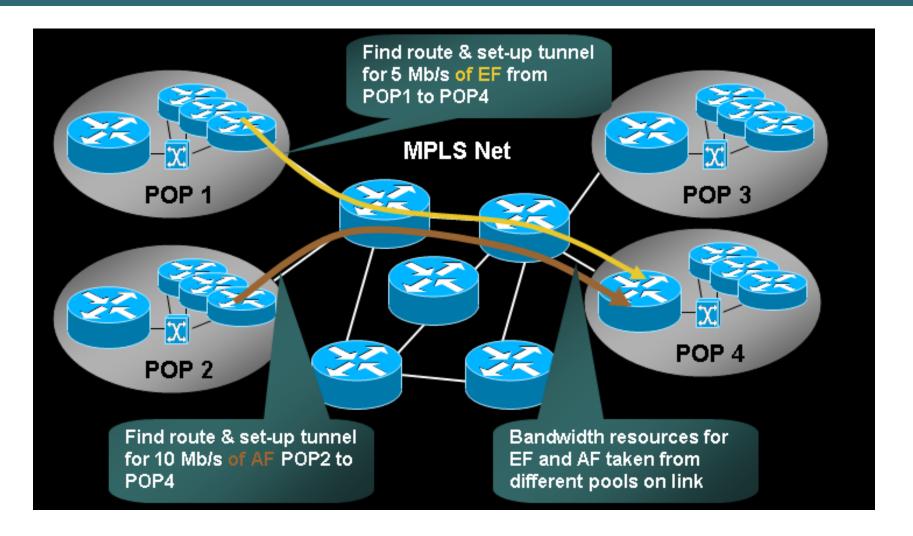
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Aggregate TE in DiffServ Network



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DiffServ TE



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DS TE and QoS

"DiffServ TE does not preclude the necessity of configuring PHB QoS in the TE path. DiffServ TE operates in conjunction with QoS mechanisms"

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Summary

- QoS techniques
 Effective allocation of network resources
- IP Diff Serv
 Service Differentiation
- MPLS & Diff Serv
 Builds scalable networks for SP
- DiffServ Tunneling Modes
 Scalable and flexible QoS options
 Supports Draft Tunneling Mode RFC
- Diff Serv TE
 Provides strict point-to-point guarantees
 Pipe Model

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Q & A